

在 Coq 中使用分离逻辑

C 语言数据结构中的单链表:

```
struct list {
    int data;
    struct list *next;
};
```

单链表的表示谓词:

```
Fixpoint sll (x: addr) (l: list Z): Assertion :=
match l with
| nil      => [] x = NULL [] && emp
| a :: l0  => [] x <> NULL [] &&
               EX y: addr,
               &(x # "list" -> "data") # Int |-> a ** 
               &(x # "list" -> "next") # Ptr |-> y **
               sll y l0
end.
```

基于 `sll` 证明 `reverse` 函数:

```
struct list *reverse(struct list *p)
/*@ With (l: list Z)
   Require sll(p, l)
   Ensure sll(_return, rev(l))
*/
{
    struct list * w = ( void * ) 0, * v = p;
    /*@ Inv exists l1 l2,
       l == app(rev(l1), l2) &&
       sll(w, l1) * sll(v, l2)
    */
    while (v) {
        struct list * t = v -> next;
        v -> next = w;
        w = v;
        v = t;
    }
    return w;
}
```

用分离逻辑描述单链表的一段:

```

Fixpoint sllseg (x y: addr) (l: list Z): Assertion :=
  match l with
  | nil      => [] x = y [] && emp
  | a :: l0 => [] x <> NULL [] &&
    EX z: addr,
    &(x # "list" -> "data") # Int |-> a ** 
    &(x # "list" -> "next") # Ptr |-> z **
    sllseg z y l0
  end.

```

`append` 的验证:

```

struct list *append(struct list *x, struct list *y)
/*@ With l1 l2
  Require sll(x, l1) * sll(y, l2)
  Ensure sll(_return, app(l1, l2))
*/
{
  struct list *t, *u;
  if (x == (void *)0) {
    return y;
  } else {
    t = x;
    u = t -> next;
    /*@ Inv
      exists l1a l1b,
      app(l1a, cons(t -> data, l1b)) == l1 &&
      t != 0 && t -> next == u &&
      sllseg(x, t, l1a) *
      sll(u, l1b) * sll(y, l2)
    */
    while (u != (void *)0) {
      t = u;
      u = t -> next;
    }
    t -> next = y;
    return x;
  }
}

```