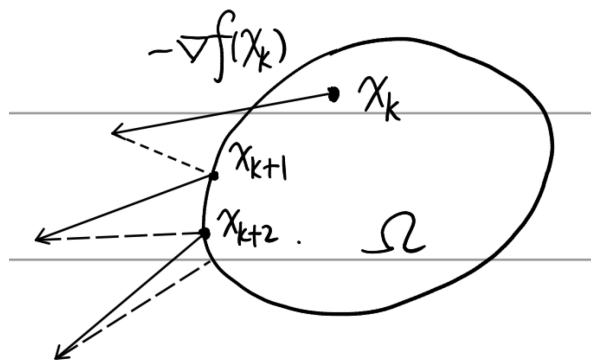


Projected Gradient Descent

1 Projection operator and projected gradient descent

To solve the inequality constrained problems, we introduce the *projected gradient descent*.

Recall the iteration step in the gradient descent method, $x_{k+1} = x_k - \eta \nabla f(x_k)$. Now we need to minimize $f(x)$ over a feasible set Ω . If $x_k - \eta \nabla f(x_k)$ is feasible, then we can run the gradient descent iteration. If $x_k - \eta \nabla f(x_k)$ is infeasible, a simple idea is to project it onto Ω . This method is called the *projected gradient descent*.



Definition (Projection)

The projection of a point onto a set is the point in the set with minimum distance to the given point. Namely, the *projection operator* is defined by

$$\mathcal{P}_{\Omega}(\mathbf{y}) = \arg \min_{\mathbf{x} \in \Omega} \|\mathbf{x} - \mathbf{y}\|.$$

The the projected gradient descent step can be given by

$$\mathbf{x}_{k+1} = \mathcal{P}_{\Omega}(\mathbf{x}_k - \eta \nabla f(\mathbf{x}_k)).$$

Let

$$\mathbf{g}(\mathbf{x}) = \frac{1}{\eta} \left(\mathbf{x} - \mathcal{P}_{\Omega}(\mathbf{x} - \eta \nabla f(\mathbf{x})) \right),$$

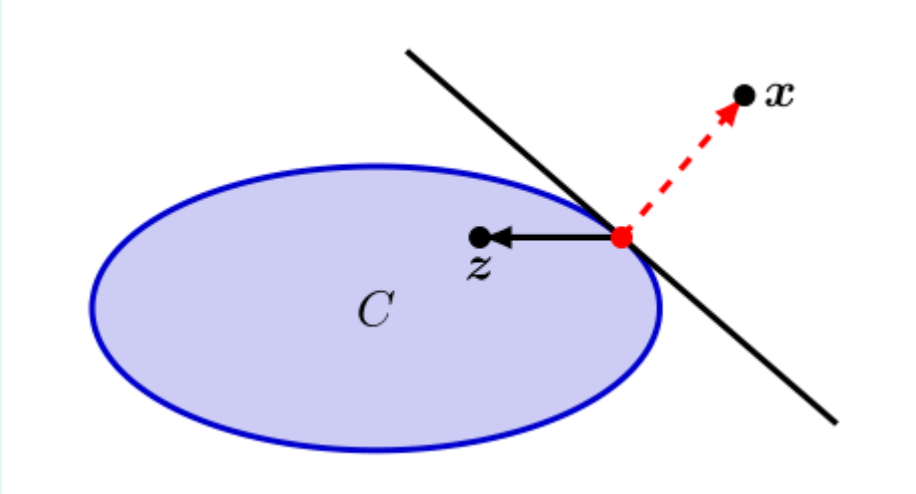
the iteration step can be expressed as

$$\mathbf{x}_{k+1} = \mathbf{x}_k - \eta \mathbf{g}(\mathbf{x}_k).$$

Recall that, in [Lecture 4](#), we show the following lemma.

Lemma

Let C be a nonempty, closed and convex set. Given $\mathbf{x}$ and $\mathbf{y} = \mathcal{P}_C(\mathbf{x})$, for any $\mathbf{z} \in C$, it holds that $\langle \mathbf{x} - \mathbf{y}, \mathbf{z} - \mathbf{y} \rangle \leq 0$.



Conversely, if there exists $\mathbf{y} \in C$ such that $\langle \mathbf{x} - \mathbf{y}, \mathbf{z} - \mathbf{y} \rangle \leq 0$, we have $\mathbf{y} = \mathcal{P}_C(\mathbf{x})$. Otherwise, let $\mathbf{w} = \mathcal{P}_C(\mathbf{x})$. Then we have

$$\langle \mathbf{x} - \mathbf{w}, \mathbf{y} - \mathbf{w} \rangle \leq 0.$$

However, we also have $\langle \mathbf{x} - \mathbf{y}, \mathbf{w} - \mathbf{y} \rangle \leq 0$, which implies that

$$\langle \mathbf{x} - \mathbf{w}, \mathbf{w} - \mathbf{y} \rangle = \langle \mathbf{x} - \mathbf{y}, \mathbf{w} - \mathbf{y} \rangle + \langle \mathbf{y} - \mathbf{w}, \mathbf{w} - \mathbf{y} \rangle < 0$$

if $\mathbf{y} \neq \mathbf{w}$. Contradiction.

Thus, $\mathbf{y} = \mathcal{P}_C(\mathbf{x})$ if and only if $\langle \mathbf{x} - \mathbf{y}, \mathbf{z} - \mathbf{y} \rangle \leq 0$ for any $\mathbf{z} \in C$.

Applying this lemma, we can show that $\mathbf{g}(\mathbf{x})$ plays a similar role as $\nabla f(\mathbf{x})$ in the gradient descent.

Lemma

For any $\mathbf{x} \in \Omega$,

$$\langle \nabla f(\mathbf{x}), \mathbf{g}(\mathbf{x}) \rangle \geq 0.$$

The inequality holds if and only if $\mathbf{g}(\mathbf{x}) = \mathbf{0}$.

Proof

Since $\mathbf{x} \in \Omega$, we have

$$\langle \mathbf{x} - \mathcal{P}_\Omega(\mathbf{x} - \eta \nabla f(\mathbf{x})), \mathbf{x} - \eta \nabla f(\mathbf{x}) - \mathcal{P}_\Omega(\mathbf{x} - \eta \nabla f(\mathbf{x})) \rangle \leq 0,$$

which gives that

$$\langle \eta \mathbf{g}(\mathbf{x}), \eta \mathbf{g}(\mathbf{x}) - \eta \nabla f(\mathbf{x}) \rangle = \eta^2 \langle \mathbf{g}(\mathbf{x}), \mathbf{g}(\mathbf{x}) - \nabla f(\mathbf{x}) \rangle \leq 0.$$

Thus,

$$\langle \nabla f(\mathbf{x}), \mathbf{g}(\mathbf{x}) \rangle \geq \langle \mathbf{g}(\mathbf{x}), \mathbf{g}(\mathbf{x}) \rangle.$$

So we know that $-\mathbf{g}(\mathbf{x})$ is a descending direction. Now we show that if $\mathbf{g}(\mathbf{x}) = \mathbf{0}$ then $\mathbf{x}$ is a minimum point.

Lemma

$\mathbf{x}^*$ is a minimum point of f over Ω , iff $\mathbf{g}(\mathbf{x}) = \mathbf{0}$, namely,
 $\mathbf{x}^* = \mathcal{P}_\Omega(\mathbf{x}^* - \eta \nabla f(\mathbf{x}^*))$.

Proof

Applying the above lemma, we have $\mathbf{x}^* = \mathcal{P}_\Omega(\mathbf{x}^* - \nabla f(\mathbf{x}^*))$ if and only if

$$\langle \mathbf{x}^* - \eta \nabla f(\mathbf{x}^*) - \mathbf{x}^*, \mathbf{z} - \mathbf{x}^* \rangle \leq 0$$

for all $\mathbf{z} \in \Omega$, which is further equivalent to

$$\langle \nabla f(\mathbf{x}^*), \mathbf{z} - \mathbf{x}^* \rangle \geq 0.$$

We conclude this lemma by the first-order optimality conditions of convex functions.

Hence, in the projected gradient descent, we can stop when $\mathbf{g}(\mathbf{x}_k)$ is small, or equivalently when $\mathbf{x}_{k+1} - \mathbf{x}_k$ is small.

Tip

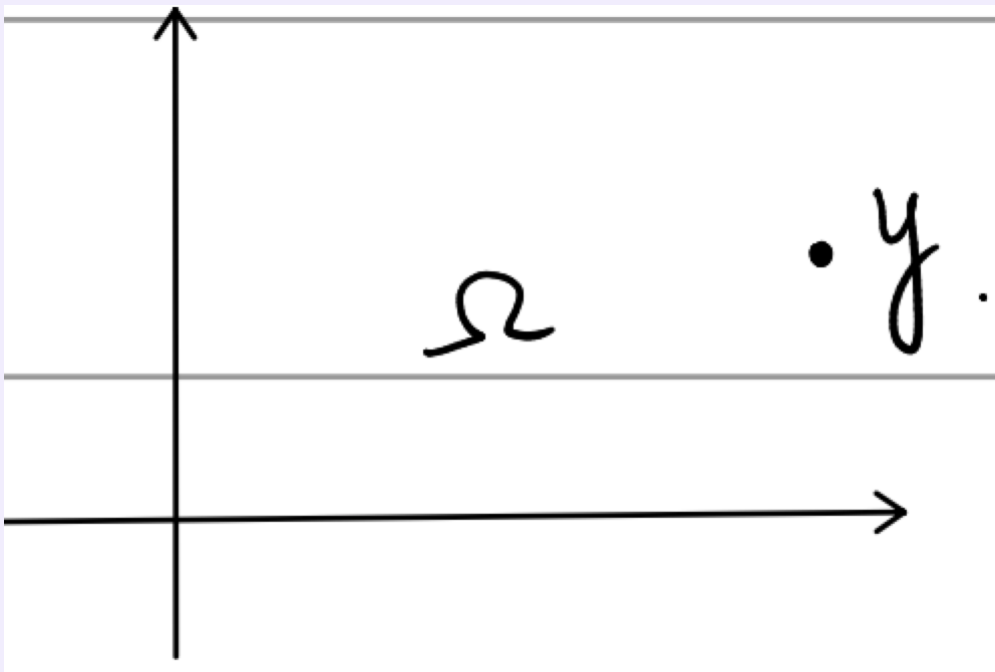
The projected gradient descent $\mathbf{x}_{k+1} = \mathcal{P}_\Omega(\mathbf{x}_k - \eta \nabla f(\mathbf{x}_k))$ can be also viewed as $\mathbf{x}_{k+1} = \arg \min_{\mathbf{z} \in \Omega} \langle \nabla f(\mathbf{x}_k), \mathbf{z} - \mathbf{x}_k \rangle + \frac{1}{2\eta} \|\mathbf{z} - \mathbf{x}_k\|^2$.

2 Examples of projection operator

Projected gradient descent is useful when the projection operator can be computed efficiently. Here we give some examples.

Example 1 (Box constraints)

$$\Omega = \{\mathbf{x} \mid a_i \leq x_i \leq b_i, \quad i = 1, \dots, n\}$$

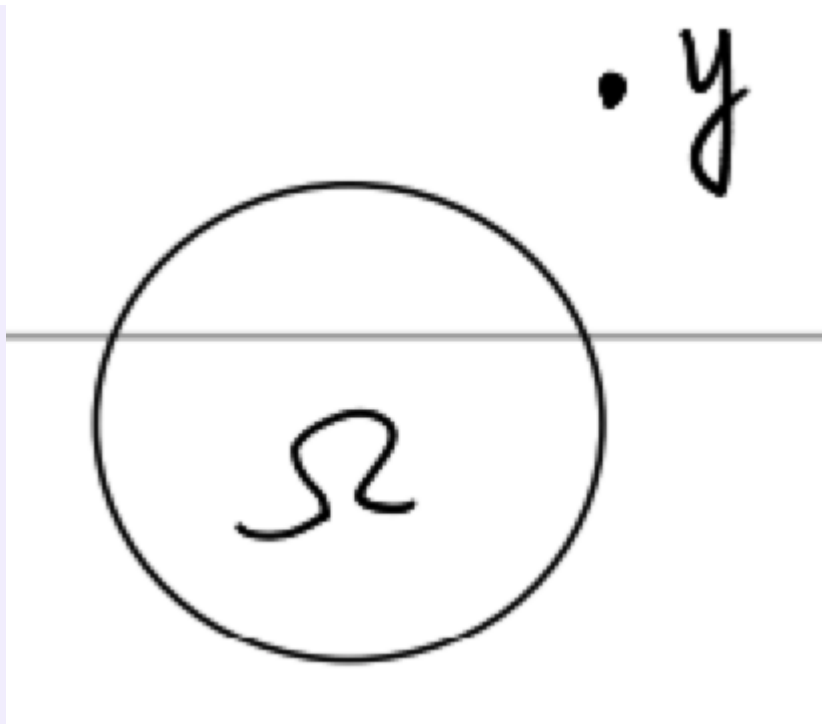


It is easy to see that

$$[\mathcal{P}_\Omega(y)]_i = \min \{b_i, \max\{a_i, y_i\}\} = \begin{cases} a_i & y_i < a_i \\ y_i & a_i \leq y_i \leq b_i \\ b_i & y_i > b_i \end{cases}$$

Example 2 (L^2 constraints, ridge regression)

$$\Omega = \{\mathbf{x} \mid \|\mathbf{x}\|_2 \leq t\}$$



The projection operator $\mathcal{P}_\Omega(y)$ is to compute

$$\begin{aligned} \min \quad & \|x - y\|^2 \\ \text{subject to} \quad & \|x\|_2^2 \leq t^2 \end{aligned}$$

By KKT condition, there exists $\mu \geq 0$ such that

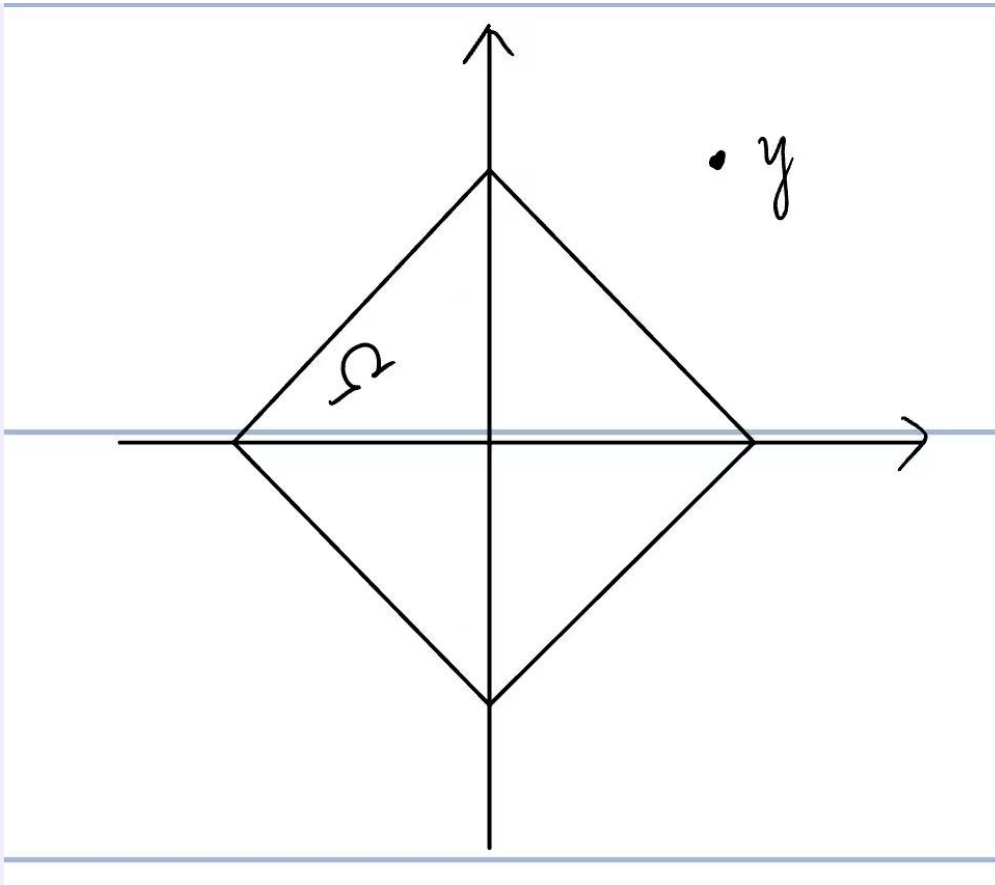
$$2(x - y) + 2\mu x = 0 \quad \text{and} \quad \mu(\|x\|^2 - t) = 0$$

Then we have $y = (1 + \mu)x$.

Hence, $\mathcal{P}_\Omega(y) = \min \left\{ 1, \frac{t}{\|y\|_2} \right\} y$.

Example 3 (L^1 constraints, LASSO)

$$\Omega = \{x : \|x\|_1 \leq t\}$$



Unfortunately, there is no closed form for the projection operator $\mathcal{P}_\Omega(y)$. But we can compute it efficiently.

By symmetry, we only need to consider the case where $y_i \geq 0$ for all i . Now $\mathcal{P}_\Omega(y)$ is equivalent to the following optimization problem:

$$\begin{aligned} \min \quad & \|x - y\|^2 \\ \text{subject to} \quad & \sum_i x_i \leq t \\ & x_i \geq 0, \forall i. \end{aligned}$$

By KKT condition, assume there exist KKT multipliers $\mu_0, \dots, \mu_n$ such that

$$\begin{cases} 2(x_i - y_i) + \mu_0 - \mu_i = 0, \forall i \\ \mu_0(\sum x_i - t) = 0 \\ \mu_i x_i = 0 \\ \mu_i \geq 0 \\ \sum x_i \leq t, x_i \geq 0 \end{cases}$$

- Case 1. $\|y\|_1 \leq t$, then $\mu_0 = \mu_i = 0$. Hence $x = y$.

- Case 2. $\|y\|_1 > t$, then

$$\sum 2(x_i - y_i) + \mu_0 - \mu_1 = 2(\sum x_i - \sum y_i) + n\mu_0 - \sum \mu_i = 0, \text{ hence } \mu_0 > 0.$$

Since $\mu_0(\sum x_i - t) = 0$, we have $\sum x_i = t$.

- If $\mu_i = 0$, by $2(x_i - y_i) + \mu_0 - \mu_i = 0$, we have $x_i = y_i - \frac{1}{2}\mu_0$.

- If $\mu_i > 0$, by $\mu_i x_i = 0$, we have $x_i = 0$.

Now we have

$$x_i = \begin{cases} y_i - \frac{1}{2}\mu_0 & \text{if } y_i \geq \frac{1}{2}\mu_0 \\ 0 & \text{otherwise} \end{cases}$$

and $\sum x_i = t$.

We may use the binary search to find μ_0 , where the lower bound is 0 and the upper bound is $\max y_i$.