

# On the Hardness of $K$ -Subspaces

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## Abstract

In this report, we study the hardness and the approximability of the  $K$ -subspace problem. The  $K$ -subspace problem is an optimization problem. It takes as input a set of  $N$  data points in  $\mathbb{R}^D$  and  $K$  positive integers  $r_1, \dots, r_K$ , and outputs  $K$  subspaces represented by  $K$  matrices  $U_1, \dots, U_K$  each of which satisfies  $U_k \in \mathbb{R}^{D \times r_k}$  and  $U_k^\top U_k = I$  such that, when clustering those  $N$  data according to the nearest subspace, the sum of the squared distances between data and their nearest subspaces is minimized. Firstly, we prove that, even for  $K = 2$  and  $r_1 = \dots = r_K$ , it is NP-complete to decide if there exists  $K$  subspaces that contains all the data, i.e., the  $K$ -subspace objective is 0. This implies that it is NP-hard to approximate the  $K$ -subspaces problem to within factor  $\alpha$ , for any finite  $\alpha > 0$ . Secondly, we prove that if  $r := r_1 = \dots = r_K$  is a fixed constant, the  $K$ -subspaces problem is APX-hard: there exists  $\varepsilon > 0$  such that it is NP-hard to approximate the  $K$ -subspaces problem to within factor  $(1 + \varepsilon)$ .

## 1 Problem Formulation

Given a dataset  $\mathcal{X} = \{x_1, \dots, x_N\} \subseteq \mathbb{R}^D$  of  $N$  data points in  $\mathbb{R}^D$  which lie near the union of  $K$  unknown subspaces  $\mathcal{S}_1, \dots, \mathcal{S}_K$  with orthonormal bases  $U_1, \dots, U_K$  and dimensions  $r_1, \dots, r_K$ , respectively. The goal of subspace clustering is to cluster the data points according to nearest subspace, which we now define formally.

**Definition 1.** The  $K$ -Subspaces (KSS) problem is an optimization problem which takes as input a data set  $\mathcal{X} = \{x_1, \dots, x_N\} \subseteq \mathbb{R}^D$  and  $K$  positive integers  $r_1, \dots, r_K$ , and outputs a partition  $\mathcal{C} = \{C_1, \dots, C_K\}$  of  $\mathcal{X}$  and a collection of  $K$  matrices  $\mathcal{U} = \{U_1, \dots, U_K \mid U_k \in \mathbb{R}^{D \times r_k}, U_k^\top U_k = I\}$  such that the objective

$$\sum_{k=1}^K \sum_{x \in C_k} \left\| x - U_k U_k^\top x \right\|_2^2 \tag{1}$$

is minimized.

We present a family of inapproximability results. Firstly, we show that even for  $K = 2$  and any finite  $\alpha > 0$  which may depend on  $D, N, r$ , approximating the KSS objective to within factor  $\alpha$  is NP-hard. Secondly, we show that the  $K$ -Subspaces is APX-hard for every fixed  $r$ . Notice that both hardness results are stronger than, and can be applied to the conventional NP-hardness result, where  $r, D, K$  and  $x_i$ 's are inputs to the  $K$ -subspaces instance.

## 2 Inapproximability for $K = 2$

We will establish a reduction from the *balanced vertex separator* problem. Definition 2 and Theorem 3 below are on a special case of the results in [4, 2].

**Definition 2.** Given a graph  $G = (V, E)$ , a *balanced vertex separator* of size  $\ell$  is a subset  $S \subseteq V$  of vertices with  $|S| = \ell$  such that  $V \setminus S$  can be partitioned into two sets  $X_1, X_2$  of equal size  $|X_1| = |X_2|$  and for any  $v_1 \in X_1, v_2 \in X_2, (v_1, v_2) \notin E$ .

The proof of Theorem 7.1 in [2] can be used to prove the theorem below. We include the proof for completeness.

**Theorem 3.** *It is NP-complete to decide if a graph  $G = (V, E)$  has a balanced vertex separator of size  $\ell$ .*

*Proof.* It is easy to see that the problem is in NP, as  $(S, X_1, X_2)$  can be used as a yes certificate.

To show the problem is NP-hard, we reduce it from a well-known problem CLIQUE: given a graph  $G = (V, E)$  and an integer  $k$ , decide if  $G$  has a clique of size  $k$ . Assume  $u_1 \in V$  is known to belong to a clique of size  $k$ , if such clique exists. Such assumption can be made without loss of generality because there are at most  $n$  possible choices for  $u_1$  and we can just try them all. We construct the balanced vertex separator instance  $(\bar{G} = (\bar{V}, \bar{E}), \ell)$  as follows. Set  $\ell = k$ .  $\bar{V}$  is defined as follows: create a vertex  $v_i$  in  $\bar{V}$  for each  $u_i \in V$ ; create a vertex  $v_{ij}$  in  $\bar{V}$  for each  $(u_i, u_j) \in E$ ; create a set of vertices  $W \subseteq \bar{V}$  with  $|W| = |V| + |E| - k^2$ .  $\bar{E}$  is defined as follows: create an edge  $(v_i, v_j) \in \bar{E}$  for each  $1 \leq i < j \leq |V|$  (i.e., the set  $V$  forms a clique in  $\bar{G}$ ); include edges  $(v_i, v_{ij}) \in \bar{E}$  and  $(v_j, v_{ij}) \in \bar{E}$  for each  $(u_i, u_j) \in E$ ; include  $(w, v_1) \in \bar{E}$  for each  $w \in W$ .

If the CLIQUE instance is a yes instance and assume without loss of generality that  $\{u_1, \dots, u_k\}$  form a clique, it is straightforward to check that  $S = \{v_1, \dots, v_k\}$ ,  $X_1 = W \cup \{v_{ij} : 1 \leq i < j \leq k\}$  and  $X_2 = \bar{V} \setminus (S \cup X_1)$  give a valid vertex separator solution: all edges going out from  $X_1$  are connected to  $S$  (in particular, none of them is connected to  $X_2$ ), and  $|X_1| = |X_2| = |V| + |E| - k - \binom{k}{2}$ .

If the CLIQUE instance is a no instance, we aim to show that  $\bar{G}$  does not have a balanced vertex separator of size  $k = \ell$ . Suppose otherwise a solution  $(S, X_1, X_2)$  exists. It must be that  $|S| = k$  and  $|X_1| = |X_2| = |V| + |E| - k - \binom{k}{2}$ . Firstly, we must include

$v_1$  in  $S$ , for otherwise the connected component including  $v_1$  after  $S$ 's removal has size at least  $|W| + |V| = 2|V| + |E| - k^2 > |V| + |E| - k - \binom{k}{2}$ , which does not fit into any of  $X_1, X_2$  (we can make sure  $|V| \gg k$  in the CLIQUE instance by creating extra isolated vertices). Next, we try to fill vertices into  $X_1$ . We can put all vertices in  $W$  into  $X_1$ , as these vertices become isolated when  $v_1$  is removed. We then still need to fill in another  $(|V| + |E| - k - \binom{k}{2}) - (|V| + |E| - k^2) = \binom{k}{2}$  vertices into  $X_1$ . The only way to do this is to choose another  $k - 1$  vertices from  $v_2, \dots, v_n$  for  $S$  such that the chosen  $k$  vertices correspond to a  $k$ -clique in  $G$ . However, this contradicts to that the CLIQUE instance is a no instance.  $\square$

Now we are ready to present our main result. The theorem below shows that it is NP-complete to decide if the minimum KSS objective value is 0, even for  $K = 2$ .

**Theorem 4.** *There exists  $r$  such that it is NP-complete to decide if there exists  $\mathcal{C} = \{c_1, c_2\}$  and  $\mathcal{U} = \{U_1, U_2\}$  such that*

$$\sum_{k=1}^2 \sum_{i: x_i \in c_k} \left\| x_i - U_k U_k^\top x_i \right\|_2^2 = 0,$$

subjecting to  $U_1, U_2 \in \mathbb{R}^{D \times r}$ .

*Proof.* It is easy to see that the problem is in NP, as we can use  $(\mathcal{C}, \mathcal{U})$  as a yes certificate.

To show the problem is NP-hard, we reduce it from the problem of finding a balanced vertex separator of size  $\ell$  in a graph  $G = (V, E)$ , whose NP-completeness is established in Theorem 3. In this proof, let  $e_i \in \mathbb{R}^{|V|}$  be the unit vector with the  $i$ -th coordinate being 1 and the remaining coordinates being 0. Given a vertex separator instance  $(G = (V, E), \ell)$ , we construct the following  $K$ -subspace instance:

- $r = \frac{1}{2}(|V| + \ell)$  (we assume  $|V| + \ell$  is even; otherwise the balanced vertex separator instance is a trivial no instance);
- We create  $N = |V| + 3|E|$  data in  $\mathbb{R}^{|V|}$  as follows.
  - For each  $v_i \in V$ , create a datum  $x_i = e_i$ ;
  - For each  $(v_i, v_j) \in E$  (assume  $i < j$ ), create three data  $y_{ij}^0 = e_i + e_j$ ,  $y_{ij}^1 = 0.5e_i + e_j$  and  $y_{ij}^2 = e_i + 0.5e_j$ .

To conclude the proof, we need to show the completeness and soundness:

1. If  $G$  has a balanced vertex separator of size  $\ell$ , there exist  $\mathcal{C}, \mathcal{U}$  such that the  $K$ -subspaces objective is 0.
2. If there exists  $\mathcal{C}, \mathcal{U}$  such that the  $K$ -subspaces objective is 0,  $G$  has a balanced vertex separator of size  $\ell$ .

To show 1, without loss of generality, let  $(S, X_1, X_2)$  be the balanced vertex separator solution with

$$X_1 = \left\{ v_1, \dots, v_{\frac{|V|-\ell}{2}} \right\}, S = \left\{ v_{\frac{|V|-\ell}{2}+1}, \dots, v_{\frac{|V|+\ell}{2}} \right\} \text{ and } X_2 = \left\{ v_{\frac{|V|+\ell}{2}+1}, \dots, v_{|V|} \right\}.$$

We consider

$$U_1 = \begin{bmatrix} e_1 & e_2 & \cdots & e_{\frac{|V|+\ell}{2}} \end{bmatrix}, \quad U_2 = \begin{bmatrix} e_{\frac{|V|-\ell}{2}+1} & e_{\frac{|V|-\ell}{2}+2} & \cdots & e_{|V|} \end{bmatrix},$$

and

$$c_1 = \left\{ x_i : 1 \leq i \leq \frac{|V|+\ell}{2} \right\} \cup \left\{ y_{ij}^0, y_{ij}^1, y_{ij}^2 : 1 \leq i < j \leq \frac{|V|+\ell}{2} \right\},$$

$$c_2 = \left\{ x_i : \frac{|V|-\ell}{2} + 1 \leq i \leq |V| \right\} \cup \left\{ y_{ij}^0, y_{ij}^1, y_{ij}^2 : \frac{|V|-\ell}{2} + 1 \leq i < j \leq |V| \right\}.$$

We can see that  $c_1 \cup c_2$  covers all the  $N$  data:  $c_1$  covers all those  $x_i$ 's corresponding to vertices in  $X_1 \cup S$ , and all those  $y_{ij}^0$ 's (or  $y_{ij}^1$ 's,  $y_{ij}^2$ 's) corresponding to edges in the subgraph induced by  $X_1 \cup S$ ;  $c_2$  covers all those  $x_i$ 's corresponding to vertices in  $S \cup X_2$ , and all those  $y_{ij}^0$ 's (or  $y_{ij}^1$ 's,  $y_{ij}^2$ 's) corresponding to edges in the subgraph induced by  $S \cup X_2$ . In particular, there is no such datum  $y_{ij}^0$  (or  $y_{ij}^1$ ,  $y_{ij}^2$ ) corresponding to an edge between  $X_1, X_2$ , as  $S$  is a vertex separator. Moreover, by our construction, it is easy to see that any  $x \in c_1$  is in the subspace spanned by all column vectors of  $U_1$ , implying that  $\|x - U_1 U_1^\top x\|_2 = 0$ . Similarly, for any  $x \in c_2$ , we have  $\|x - U_2 U_2^\top x\|_2 = 0$ . This implies that the  $K$ -subspaces objective is 0, concluding 1.

To show 2, suppose there exist  $\mathcal{C}, \mathcal{U}$  such that the  $K$ -subspaces objective is 0. First of all, for each pair  $i, j$  with  $(v_i, v_j) \in E$ , we can assume without loss of generality that the three data  $y_{ij}^0, y_{ij}^1, y_{ij}^2$  belong to the same cluster (either  $c_1$  or  $c_2$ ). In fact,  $y_{ij}^0, y_{ij}^1, y_{ij}^2$  are linearly dependent. If any two of them are in the same subspace defined by either  $U_1$  or  $U_2$ , the third one is also in this subspace. Thus,  $y_{ij}^0, y_{ij}^1, y_{ij}^2$  are either in three different subspaces, or in a single subspace. Since we are only allow  $K = 2$  subspaces, the  $K$ -subspaces objective being 0 implies that they are in a single subspace.

Let  $V_1 \subseteq V$  be defined as

$$V_1 = \{v_i \in V : x_i \in c_1\} \cup \{v_i \in V : \exists j \text{ such that either } y_{ij}^0 \in c_1 \text{ or } y_{ji}^0 \in c_1\}.$$

In words,  $V_1$  is the set of vertices  $v_i$  such that either the datum corresponding to  $v_i$  is in  $c_1$  or the datum corresponding to an edge incident to  $v_i$  is contained in  $c_1$ . Similarly, define  $V_2 \subseteq V$  as

$$V_2 = \{v_i \in V : x_i \in c_2\} \cup \{v_i \in V : \exists j \text{ such that either } y_{ij}^0 \in c_2 \text{ or } y_{ji}^0 \in c_2\}.$$

We will show that the subspace defined by  $U_1$  has dimension at least  $|V_1|$ . To see this, consider the matrix  $M_1 \in \mathbb{R}^{|V| \times |c_1|}$  whose columns are those data from  $c_1$ . By our construction,  $M_1$  has exactly  $|V_1|$  rows that are not all-zero. We claim that these  $|V_1|$  rows are

linearly independent. Suppose this is not true and there exists coefficients  $a_1, \dots, a_{|V_1|}$  such that the weighted sum of these  $|V_1|$  rows is the zero vector and at least one of  $a_1, \dots, a_{|V_1|}$  is nonzero. Suppose  $a_1 \neq 0$  and the  $i$ -th row of  $M_1$  corresponds to  $a_1$ . Consider an arbitrary nonzero entry in the  $i$ -th row. The column containing this entry either corresponds to the datum  $x_i$  or one of  $y_{ij}^0, y_{ij}^1, y_{ij}^2$ . In the former case, this entry is the only nonzero entry this column, and assigning  $a_1 \neq 0$  to row  $i$  cannot make the weighted sum of those  $|V_1|$  rows a zero vector, which is a contradiction. In the latter case, the  $|V| \times 3$  submatrix defined by the three columns representing  $y_{ij}^0, y_{ij}^1, y_{ij}^2$  ( $M_1$  contains all these three columns, because  $y_{ij}^0, y_{ij}^1, y_{ij}^2$  belongs to the same cluster as we have shown earlier) is

$$\begin{bmatrix} 0 & 0 & 0 \\ \vdots & \vdots & \vdots \\ 0 & 0 & 0 \\ 1 & 0.5 & 1 \\ 0 & 0 & 0 \\ \vdots & \vdots & \vdots \\ 0 & 0 & 0 \\ 1 & 1 & 0.5 \\ 0 & 0 & 0 \\ \vdots & \vdots & \vdots \\ 0 & 0 & 0 \end{bmatrix},$$

where row  $i$  and  $j$  are the only two rows that are not all zero. If  $a_1 \neq 0$  is assigned to row  $i$ , it is impossible to assign a coefficient to row  $j$  such that the weighted sum of the two rows is a zero vector. This again is a contradiction. Therefore, the  $|V_1|$  nonzero rows in  $M_1$  are linearly independent, so the rank of  $M_1$  is  $|V_1|$ . Viewing this rank as column rank, this implies that the subspace containing all the data in  $c_1$  has dimension at least  $|V_1|$ , so  $U_1$  has dimension at least  $|V_1|$ . Thus,  $|V_1| \leq \dim(U_1) = r = \frac{1}{2}(|V| + \ell)$ . Similarly,  $|V_2| \leq \frac{1}{2}(|V| + \ell)$ . We arbitrarily add extra vertices in  $V_1, V_2$  such that  $|V_1| = |V_2| = \frac{1}{2}(|V| + \ell)$ .

Since  $c_1 \cup c_2$  contains all the data including  $x_1, \dots, x_{|V|}$ , we have  $V_1 \cup V_2 = V$ . By inclusion-exclusion principle,  $|V_1 \cap V_2| = \frac{1}{2}(|V| + \ell) + \frac{1}{2}(|V| + \ell) - |V| = \ell$ . We claim that  $S = V_1 \cap V_2$ ,  $X_1 = V_1 \setminus V_2$  and  $X_2 = V_2 \setminus V_1$  give a valid solution to the balanced vertex separator instance. It suffices to show that  $(v_1, v_2) \notin E$  for any  $v_1 \in X_1 = V_1 \setminus V_2$  and  $v_2 \in X_2 = V_2 \setminus V_1$ . Suppose this is not the case and  $(v_1, v_2) \in E$ . There is a datum  $y_{12}^0$  and it must be contained in either  $c_1$  or  $c_2$ . Assume without loss of generality that  $y_{12}^0 \in c_1$ . By our construction, both endpoints  $v_1, v_2$  should be contained in  $V_1$ , which contradicts to  $v_2 \in V_2 \setminus V_1$ . We conclude 2 here.  $\square$

Theorem 4 straightforwardly implies the following corollary, which shows that the  $K$ -subspaces problem is totally inapproximable.

**Corollary 5.** For any finite  $\alpha > 0$  which may depend on  $D, N, K, r$ , approximating the  $K$ -subspaces objective function to within factor  $\alpha$  is NP-hard, even for  $K = 2$ .

### 3 Inapproximability for Fixed $r$

#### 3.1 Inapproximability Result for $r = 1$

We first prove the following lemma.

**Lemma 6.** Let  $t, r$  be two positive real numbers with  $t > r$ . Let  $x \in \mathbb{R}^d$  with  $\|x\|_2 \leq r$  and  $y = (x, t)^\top \in \mathbb{R}^{d+1}$ , and  $\mathcal{S} = \{(x_1, \dots, x_d, x_{d+1}) \in \mathbb{R}^{d+1} : x_{d+1} = t, \sum_{i=1}^d x_i^2 \leq r^2\}$  be the  $d$ -dimensional hyper-ball in the affine plane  $\mathcal{A} = \{(x_1, \dots, x_d, x_{d+1}) \in \mathbb{R}^{d+1} : x_{d+1} = t\}$ , with center  $p = (0, \dots, 0, t)$  and radius  $r$ . Let  $\mu = (\mu_1, \dots, \mu_{d+1}) \in \mathbb{R}^{d+1}$  be a unit vector. Let  $\nu = \frac{t}{\mu_{d+1}}(\mu_1, \dots, \mu_d) \in \mathbb{R}^d$ . We have

$$\|y - \mu\mu^\top y\|_2 \leq \|x - \nu\|_2.$$

If in addition the line containing  $\mu$  intersects  $\mathcal{S}$ , we also have

$$\|x - \nu\|_2 \leq \frac{t^2 + r^2}{t^2 - r^2} \|y - \mu\mu^\top y\|_2. \quad (2)$$

*Proof.* Firstly, it is straightforward to show that  $\|x - \nu\|_2 = \|y - \frac{t}{\mu_{d+1}}\mu\|_2$ .

Next, elementary calculus shows that  $\xi = \mu^\top y$  is the minimizer for the function  $f(\xi) := \|y - \xi\mu\|_2^2$ , so we have  $\|y - \mu\mu^\top y\|_2 \leq \|y - \frac{t}{\mu_{d+1}}\mu\|_2$ , implying the first inequality in the lemma.

For the inequality in Eq. (2), let  $u = \frac{t}{\mu_{d+1}}\mu$ . We need to show that

$$\frac{\|y - \mu\mu^\top y\|_2}{\|x - \nu\|_2} = \frac{\|y - \mu\mu^\top y\|_2}{\|y - u\|_2} \geq \frac{t^2 - r^2}{t^2 + r^2}.$$

The vector  $\mu\mu^\top y$  is the projection of  $y$  onto the subspace spanned by  $u$ . Also let  $\mathcal{H} = \{x : x = u + \alpha(y - u), \alpha \in \mathbb{R}\}$  (in words,  $\mathcal{H}$  is the affine line passing through both  $y$  and  $u$ ), and let  $q$  be the projection of the zero vector  $\underline{0}$  onto  $\mathcal{H}$ . Note that  $q, y, u$  are colinear, as are  $u, \mu\mu^\top y, \underline{0}$ . See Fig. 1 for an illustration. We have

$$\frac{\|y - \mu\mu^\top y\|_2}{\|y - u\|_2} = \frac{\|q\|_2}{\|u\|_2} = \frac{q^\top u}{\|q\|_2 \|u\|_2}, \quad (3)$$

where the first equality is based on that the two right triangles in Fig. 1 sharing the vertex  $u$  are similar, and the second equality is by the definition of projection (the projection of  $u$  onto the line passing through  $\underline{0}$  and  $q$  is exactly  $q$ , so  $q^\top u = u^\top q = \|q\|_2 \cdot \|q\|_2$ ).

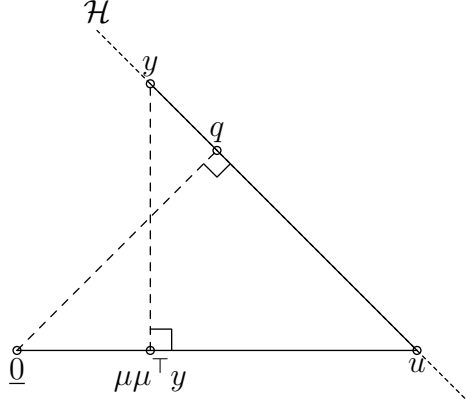


Figure 1: Illustration for Eq. (3).

Next, we will now show that  $q, y, u \in \mathcal{S}$ . By our assumption that the line containing  $\mu$  intersects  $\mathcal{S}$  and  $u$  is precisely this intersection, we have  $u \in \mathcal{S}$ . Since  $p$  is the center of  $\mathcal{S}$  and  $\|y - p\|_2 = \|x\|_2 \leq r$ , we also have  $y \in \mathcal{S}$ . To show  $q \in \mathcal{S}$ , we show that  $q$  is the projection of  $p$  (the center of  $\mathcal{S}$ ) onto  $\mathcal{H}$ . Suppose otherwise that  $\hat{q} \in \mathcal{H}$  is the projection of  $p$  onto  $\mathcal{H}$  and  $\hat{q} \neq q$ , then  $\|p - \hat{q}\|_2 < \|p - q\|_2$ . Since for all  $a \in \mathcal{A}$ ,  $(a - p)^\top p = 0$ , and  $q, \hat{q} \in \mathcal{H} \subseteq \mathcal{A}$ , we have that

$$\begin{aligned} \|\underline{0} - \hat{q}\|_2^2 &= \|\underline{0} - p\|_2^2 + \|p - \hat{q}\|_2^2 && \text{(Pythagorean theorem)} \\ &< \|\underline{0} - p\|_2^2 + \|p - q\|_2^2 \\ &= \|\underline{0} - q\|_2^2, && \text{(Pythagorean theorem)} \end{aligned}$$

giving us a contradiction since  $q$ , instead of  $\hat{q}$ , was defined as the projection of  $\underline{0}$  onto  $\mathcal{H}$ . Thus,  $q$  is the projection of  $p$  onto  $\mathcal{H}$ , implying  $\|p - q\|_2 \leq \|p - w\|_2$  for all  $w \in \mathcal{H}$ , including  $w = y$ . Therefore,  $\|p - q\|_2 \leq \|p - y\|_2 = \|x\|_2 \leq r$ , implying  $q \in \mathcal{S}$ .

Finally, we will bound  $q^\top u / \|q\|_2 \|u\|_2$  on the right hand side of Eq. (3). Since  $q, u \in \mathcal{S}$  we have by the Pythagorean theorem that  $\|q\|_2^2 \leq \|p\|_2^2 + r^2 = t^2 + r^2$  and similarly  $\|u\|_2^2 \leq t^2 + r^2$ . Thus, we have the following bound for the denominator.

$$\|q\|_2 \|u\|_2 \leq \sqrt{t^2 + r^2} \cdot \sqrt{t^2 + r^2} = t^2 + r^2.$$

For the numerator, we note that

$$q^\top u = t^2 + \sum_{i=1}^d q_i u_i \geq t^2 - \frac{1}{2} \left( \sum_{i=1}^d q_i^2 + \sum_{i=1}^d u_i^2 \right) \geq t^2 - r^2,$$

where the second last inequality is based on the inequality  $ab \geq -\frac{1}{2}(a^2 + b^2)$  and the last inequality is based on that  $q, u \in \mathcal{S}$ .

Putting these together, we have

$$\frac{q^\top u}{\|q\|_2 \|u\|_2} \geq \frac{t^2 - r^2}{t^2 + r^2}. \quad (4)$$

Eq. (3) and Eq. (4) complete the proof.  $\square$

Now we are ready to present our hardness result. The theorem below shows that the  $K$ -subspace problem is APX-hard even for  $r = 1$ .

**Theorem 7.** *For  $r = 1$ , there exists  $\varepsilon > 0$  such that it is NP-hard to approximate the optimal KSS objective*

$$\min_{\mathcal{C}, \mathcal{U}} \sum_{k=1}^K \sum_{i: x_i \in c_k} \left\| x_i - U_k U_k^\top x_i \right\|_2^2$$

to within a factor of  $(1 + \varepsilon)$ , subjecting to  $U_k \in \mathbb{R}^{D \times r}$  for each  $k = 1, \dots, K$ .

*Proof.* We will prove the lemma by a reduction from the Euclidean  $K$ -means problem. In [1], it has been shown that *there exists a constant  $\epsilon > 0$  such that it is NP-hard to approximate the  $K$ -means objective to within a factor of  $(1 + \epsilon)$ .*

Given a  $K$ -means instance  $(\{\bar{x}_1, \dots, \bar{x}_N : \bar{x}_i \in \mathbb{R}^{\bar{D}}\}, \bar{K})$ , we construct the following KSS instance:

$$\left( \{x_1, \dots, x_N : x_i \in \mathbb{R}^D\}, K \right) \quad \text{with} \quad \begin{aligned} &K = \bar{K}, D = \bar{D} + 1, \text{ and} \\ &x_i = \begin{bmatrix} \bar{x}_i \\ T \end{bmatrix} \text{ for each } i = 1, \dots, N, \end{aligned}$$

where  $T > 0$  is a very large real number such that  $\left(\frac{T^2 - 9d^2}{T^2 + 9d^2}\right)^2 (1 + \epsilon) - 1 > 0$  with  $d = \max_i \|\bar{x}_i\|_2$ . To prove the lemma, set

$$\varepsilon = \left(\frac{T^2 - 9d^2}{T^2 + 9d^2}\right)^2 (1 + \epsilon) - 1,$$

and we aim to show that a  $(1 + \varepsilon)$ -approximation algorithm for the above KSS instance will give a  $(1 + \epsilon)$ -approximation to the optimal objective value of the  $K$ -means instance.

For any solution  $(\mathcal{C} = \{c_1, \dots, c_K\}, \mathcal{U} = \{U_1, \dots, U_K\})$  to the KSS instance, let

$$\phi_{KSS}(\mathcal{C}, \mathcal{U}) = \sum_{k=1}^K \sum_{i: x_i \in c_k} \left\| x_i - U_k U_k^\top x_i \right\|_2^2$$

be the corresponding KSS objective value. Similarly, for any solution  $(\mathcal{C} = \{c_1, \dots, c_K\}, \mathcal{S} = \{s_1, \dots, s_K\})$  to the  $K$ -means instance where  $s_1, \dots, s_K$  are the  $K$  centers, let

$$\phi_{KM}(\mathcal{C}, \mathcal{S}) = \sum_{k=1}^K \sum_{i: \bar{x}_i \in c_k} \|\bar{x}_i - s_k\|_2^2$$



be the corresponding  $K$ -means objective value. Let  $(\mathcal{C}^* = \{c_1^*, \dots, c_K^*\}, \mathcal{U}^* = \{U_1^*, \dots, U_K^*\})$  be the optimal solution to the KSS instance and  $(\mathcal{C}^\dagger = \{c_1^\dagger, \dots, c_K^\dagger\}, S^\dagger = \{s_1^\dagger, \dots, s_K^\dagger\})$  be the optimal solution to the  $K$ -means instance.

Suppose we have the said  $(1 + \varepsilon)$ -approximation algorithm for KSS, and obtain the approximately optimal solution  $(\mathcal{C} = \{c_1, \dots, c_K\}, \mathcal{U} = \{U_1, \dots, U_K\})$  to the KSS instance:

$$\phi_{KSS}(\mathcal{C}, \mathcal{U}) \leq (1 + \varepsilon)\phi_{KSS}(\mathcal{C}^*, \mathcal{U}^*). \quad (5)$$

For each  $k = 1, \dots, K$ , let

$$U_k = [u_{k1} \quad u_{k2} \quad \cdots \quad u_{k\bar{D}} \quad u_{kD}]^\top,$$

and construct

$$s_k = \frac{T}{u_{kD}} \cdot [u_{k1} \quad u_{k2} \quad \cdots \quad u_{k\bar{D}}]^\top.$$

We will show that the obtained  $K$ -means solution  $(\mathcal{C} = \{c_1, \dots, c_K\}, S = \{s_1, \dots, s_K\})$  is a  $(1 + \varepsilon)$ -approximation:

$$\phi_{KM}(\mathcal{C}, S) \leq (1 + \varepsilon)\phi_{KM}(\mathcal{C}^\dagger, S^\dagger). \quad (6)$$

Before showing (6), we construct the following solution to the KSS instance based on the  $K$ -means optimal solution  $(\mathcal{C}^\dagger, S^\dagger)$ :

$$(\mathcal{C}^\dagger, \mathcal{U}' = \{U'_1, \dots, U'_K\}),$$

where

- the constructed KSS solution has the same partition, and
- $U'_k = \frac{(s_k^\dagger, T)}{\|(s_k^\dagger, T)\|_2}$ . That is, each center point  $s_k^\dagger \in \mathbb{R}^{\bar{D}}$  is appended with an extra coordinate  $T$ , and the resultant  $D$ -dimensional vector is then normalized.

The remaining part of this proof aims to show (6), and we show it by showing the following three inequalities.

$$\phi_{KM}(\mathcal{C}, S) \leq \left(\frac{T^2 + 9d^2}{T^2 - 9d^2}\right)^2 \phi_{KSS}(\mathcal{C}, \mathcal{U}) \quad (7)$$

$$\phi_{KSS}(\mathcal{C}^*, \mathcal{U}^*) \leq \phi_{KSS}(\mathcal{C}^\dagger, \mathcal{U}') \quad (8)$$

$$\phi_{KSS}(\mathcal{C}^\dagger, \mathcal{U}') \leq \phi_{KM}(\mathcal{C}^\dagger, S^\dagger) \quad (9)$$

Let  $\mathcal{S} = \{(x_1, \dots, x_D) \in \mathbb{R}^D : x_D = T, \sum_{i=1}^{\bar{D}} x_i \leq (3d)^2\}$  be the  $\bar{D}$ -dimensional ball in the affine plane  $\{(x_1, \dots, x_D) \in \mathbb{R}^D : x_D = T\}$  with center  $(0, \dots, 0, T)$  and radius  $3d$ . By

our definition of  $d$ , all the data  $x_i$ 's are within the ball with center  $(0, \dots, 0, T)$  and radius  $d$ . Loosely speaking, we can assume without loss of generality that the line containing each  $U_k$  intersects  $\mathcal{S}$ : if any  $U_k$  misses  $\mathcal{S}$ , it is too far away from all the data  $x_i$ 's, and straightforward geometric arguments can show that it is even better to replace  $U_k$  by the vector  $(0, \dots, 0, 1)$  that pointing towards the center of  $\mathcal{S}$ ; thus, upon receiving the output  $(\mathcal{C}, \mathcal{U})$  of the mentioned  $(1 + \varepsilon)$ -approximation algorithm, we can assume there is an extra cleanup step which replaces any  $U_k$  that misses  $\mathcal{S}$  by  $(0, \dots, 0, 1)$ .

By the second inequality in Lemma 6,

$$\begin{aligned} \phi_{KM}(\mathcal{C}, \mathcal{S}) &= \sum_{k=1}^K \sum_{i: \bar{x}_i \in c_k} \|\bar{x}_i - s_k\|_2^2 \leq \sum_{k=1}^K \sum_{i: x_i \in c_k} \left( \frac{T^2 + 9d^2}{T^2 - 9d^2} \right)^2 \left\| x_i - U_k U_k^\top x_i \right\|_2^2 \\ &= \left( \frac{T^2 + 9d^2}{T^2 - 9d^2} \right)^2 \phi_{KSS}(\mathcal{C}, \mathcal{U}), \end{aligned}$$

which concludes Eq. (7).

Since  $(\mathcal{C}^*, \mathcal{U}^*)$  is the optimal solution to the KSS instance while  $(\mathcal{C}^\dagger, \mathcal{U}')$  may not be optimal, (8) holds. A straightforward application of the first inequality of Lemma 6 can also show (9).

Finally, applying inequalities (7), (5), (8) and (9) one-by-one, and by our definition of  $\varepsilon$ , we obtain Eq. (6).  $\square$

### 3.2 Inapproximability Result for General Fixed $r$

In this section, we prove the following theorem.

**Theorem 8.** *For any fixed constant  $r$ , there exists  $\varepsilon > 0$  such that it is NP-hard to approximate the optimal KSS objective*

$$\min_{\mathcal{C}, \mathcal{U}} \sum_{k=1}^K \sum_{i: x_i \in c_k} \left\| x_i - U_k U_k^\top x_i \right\|_2^2$$

*to within a factor of  $(1 + \varepsilon)$ , subjecting to  $U_k \in \mathbb{R}^{D \times r}$  for each  $k = 1, \dots, K$ .*

Since the theorem for  $r = 1$  has been proved in the last section, we consider  $r \geq 2$  in this section.

**The construction** Given a  $K$ -means instance  $(\{\bar{x}_1, \dots, \bar{x}_{\bar{N}} : \bar{x}_i \in \mathbb{R}^{\bar{D}}\}, \bar{K})$ , we construct the KSS instance as follows.

- $K = \bar{K}$  and  $D = \bar{D} + r$ ;

- For each  $\bar{x}_n$  in the  $K$ -means instance, construct  $M^{r-1}$  data for the KSS instance

$$X_n = \{x_n^{i_1, \dots, i_{r-1}} : i_1, \dots, i_{r-1} \in \{0, 1, \dots, M-1\}\}$$

such that

$$x_n^{i_1, \dots, i_{r-1}} = [\bar{x}_n \quad T \quad i_1 \cdot S \quad \cdots \quad i_{r-1} \cdot S]^\top \in \mathbb{R}^D,$$

where  $T, S > 0$  are two large real numbers and  $M \in \mathbb{Z}^+$  is a large integer to be defined later.

Notice that the total number of data in the constructed KSS instance is  $N = M^{r-1} \cdot \bar{N}$ .

Let  $d = \max_{i \in \{1, \dots, \bar{N}\}} \|\bar{x}_i\|$ . For each  $i_1, \dots, i_{r-1}$ , let

$$\mathcal{S}^{i_1, \dots, i_{r-1}} = \left\{ x = (x_1, \dots, x_D) \in \mathbb{R}^D : x_{\bar{D}+1} = T, x_{\bar{D}+2} = i_1 S, \dots, x_{\bar{D}+r} = i_{r-1} S, \sum_{i=1}^{\bar{D}} x_i \leq d^2 \right\}$$

be the hyper-ball in the affine hyper-plane

$$\{x = (x_1, \dots, x_D) \in \mathbb{R}^D : x_{\bar{D}+1} = T, x_{\bar{D}+2} = i_1 S, \dots, x_{\bar{D}+r} = i_{r-1} S\}$$

centered at  $(0, \dots, 0, T, i_1 S, \dots, i_{r-1} S)$  with radius  $d$ . Although we will define  $T, S, M$  later, it is helpful to first notice that we will make  $S \gg 2d$  so that all those hyper-balls are far away from each other. Let

$$\mathcal{B} = \left\{ x = (x_1, \dots, x_D) \in \mathbb{R}^D : x_{\bar{D}+1} = T, \sum_{i \neq \bar{D}+1} x_i^2 \leq (\sqrt{r}MS + d)^2 \right\}$$

be the hyper-ball that is large enough to contain all the data points in our KSS instance. We will, in addition, make  $T$  significantly larger than the radius  $\sqrt{r}MS + d$  of  $\mathcal{B}$ , such that the line passing through the origin and an arbitrary point in  $\mathcal{B}$  has very small angle to the  $(\bar{D} + 1)$ -th unit vector.

**The high level ideas of the proof** The  $K$ -means instance is duplicated  $M^{r-1}$  times such that each  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  contains a copy of the instance. Moreover, the centers of the hyper-balls  $\{\mathcal{S}^{i_1, \dots, i_{r-1}}\}$  form a lattice structure with spacing  $S$ , and the data corresponding to each  $K$ -means instance copy are very close (compared to the spacing  $S$ ) to the respective center.

In an (approximately) optimal solution to the constructed KSS instance, each subspace  $U_k$  needs to intersect, or at least not far away from, all the  $M^{r-1}$  hyper-balls. Specifically, let  $e_i$  be the  $i$ -th unit vector in  $\mathbb{R}^D$ , the  $r$  columns of  $U_k$  need to be approximately  $e_{\bar{D}+1}, \dots, e_{\bar{D}+r}$ . Otherwise, if the  $(j+1)$ -th column has a large difference to  $e_{\bar{D}+j+1}$ , then, for any  $i_1, \dots, i_{j-1}, i_{j+1}, \dots, i_{r-1}$ ,  $U_k$  will intersect very few hyper-balls from

$\{\mathcal{S}^{i_1, \dots, i_{r-1}} : r_j = 0, 1, \dots, M-1\}$ . The total number of hyper-balls that  $U_k$  intersects is at most about  $c \cdot M^{r-2} \ll M^{r-1}$  (for some  $c \ll M$ ), and the cost for allocating those data in the not-intersected balls in the cluster  $C_k$  is high. Therefore, to avoid the high cost, the data in each ball that  $U_k$  does not intersect can only be allocated to one of the remaining  $K-1$  subspaces that intersect (or is close to) the ball. As a result, although such  $U_k$  may do a better job for those data in the  $cM^{r-2}$  balls that it intersects, but it will leave the data in the remaining  $M^{r-1} - cM^{r-2}$  balls to the remaining  $K-1$  subspaces to cover. Since  $M^{r-1} - cM^{r-2}$  is considerably larger than  $cM^{r-2}$ , switching  $U_k$  to some  $U'_k$  that intersects all the balls is much more beneficial.

Next, if each  $U_k$  intersects all the  $M^{r-1}$  balls, it is easy to see that the intersection between  $U_k$  and each ball can only be a single point: intuitively, if  $U_k$  spends more than 0 dimension to a ball, it will not have enough dimensions to cover all the balls whose centers form an  $(r-1)$ -dimensional lattice. By showing a lemma that is similar to Lemma 6, we can show that the distance between a data point and its closest subspace  $U_k$  is approximately the distance between it and the point at which  $U_k$  and the ball containing the data point intersect. For each ball  $\mathcal{S}^{i_1, \dots, i_{r-1}}$ , by viewing its  $K$  intersections to  $U_1, \dots, U_K$  as the  $K$  centers in the  $K$ -means problem, we find an approximate solution to the  $K$ -means instance. If we have an  $\varepsilon$ -approximated optimal solution to the KSS instance, we have  $M^{r-1}$  solutions to the  $K$ -means instance. Taking the best solution from them gives us an approximated optimal solution to the  $K$ -means instance.

**The proof** Let  $d$  and  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  be as defined in the last part. Let

$$\mathcal{P}^{i_1, \dots, i_{r-1}} = \{(x_1, \dots, x_D) : x_{\bar{D}+1} = T, x_{\bar{D}+2} = i_1 S, \dots, x_D = i_{r-1} S\}$$

be the hyperplane that contains  $\mathcal{S}^{i_1, \dots, i_{r-1}}$ . We will prove Theorem 8 following the ideas in the last part.

**Proposition 9.** *If a subspace  $U_k \in \mathbb{R}^{D \times r}$  intersects at least  $M^{r-2} + 1$  hyperplanes from  $\{\mathcal{P}^{i_1, \dots, i_{r-1}} : i_1, \dots, i_{r-1} \in \{0, 1, \dots, M-1\}\}$ , then the intersection between  $U_k$  and a hyperplane it intersects is a single point. In particular, if a subspace  $U_k \in \mathbb{R}^{D \times r}$  intersects at least  $M^{r-2} + 1$  hyper-balls from  $\{\mathcal{S}^{i_1, \dots, i_{r-1}} : i_1, \dots, i_{r-1} \in \{0, 1, \dots, M-1\}\}$ , then the intersection between  $U_k$  and a hyper-ball it intersects is a single point.*

*Proof.* For the ease of notation, we index each hyperplane  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  by a single integer  $j \in \{0, 1, \dots, M^{r-1} - 1\}$ :

$$\mathcal{P}_j := \mathcal{P}^{i_1, \dots, i_{r-1}} \quad \text{where } j = \sum_{s=1}^{r-1} i_s \cdot M^{s-1}.$$

Suppose for the sake of contradiction that  $U_k$  intersects at least  $M^{r-2} + 1$  hyper-balls and there exists  $\mathcal{P}_{j^*}$  such that  $U_k$  intersects it at at least two points:  $z_{j^*}$  and  $z'_{j^*}$ . For

each  $j$  such that  $U_k$  intersects  $\mathcal{P}_j$ , let  $z_j$  be the intersection of  $U_k$  and  $\mathcal{P}_j$  (or one of the intersection points if the intersection is more than a single point). For each  $z_j$ , let  $z_j^- \in \mathbb{R}^r$  be the truncated vector which consists of the last  $r$  entries of  $z_j$ . That is, the first entry of  $z_j^-$  is  $T$ , and the remaining  $r - 1$  entries are the values of  $i_1 S, \dots, i_{r-1} S$ . As a first step, we show that there exists  $r - 1$  indices  $j_1, \dots, j_{r-1}$  such that  $U_k$  intersects  $\mathcal{P}_{j_1}, \dots, \mathcal{P}_{j_{r-1}}$  at  $z_{j_1}, \dots, z_{j_{r-1}}$  and the  $r$  vectors  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r-1}}^-$  are linearly independent.

Suppose this is not the case, and we can find at most  $r' \leq r - 2$  indices  $j_1, \dots, j_{r'}$  satisfying this. For any  $j$  such that  $U_k$  intersects  $\mathcal{P}_j$  at  $z_j$ , the corresponding  $z_j^-$  must be in the  $r'$ -dimensional subspace spanned by  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r'}}^-$ . Let  $Z = \{(z_1, \dots, z_r) : z_1 = T; z_2, \dots, z_r \in \{0, 1, \dots, M - 1\}\}$  be the set of all possible values of  $z_j^-$ . We will derive a contradiction by showing that the subspace spanned by  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r'}}^-$  overlaps to  $Z$  at at most  $M^{r'}$  points (since  $M^{r'} \leq M^{r-2}$ , this contradicts to that  $U_k$  intersects at least  $M^{r-2} + 1$  hyper-balls). Since  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r'}}^-$  are linearly independent, we can find  $t_1, \dots, t_{r'+1}$  such that the  $(r' + 1) \times (r' + 1)$  submatrix from the rank- $(r' + 1)$  matrix  $[z_{j^*}^- \ z_{j_1}^- \ \dots \ z_{j_{r'}}^-]$  obtained by taking rows  $t_1, \dots, t_{r'+1}$  is non-degenerated. Moreover, we choose  $t_1 = 1$ . Then, for each  $z_j^-$  that is given by a linear combination of  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r'}}^-$ , once the  $t_1$ -th, ...,  $t_{r'+1}$ -th entries are determined, the remaining  $r - r' - 1$  entries are also determined. Since the  $t_1$ -th entry, which is the first entry, is always  $T$ , the number of possible choices for the  $r' + 1$  entries, the  $t_1$ -th, ...,  $t_{r'+1}$ -th entries, are  $M^{r'}$ . This implies that there are at most  $M^{r'}$  points from  $Z$  that are in the subspace spanned by  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r'}}^-$ , which gives us the desired contradiction.

In the next step, we show that, if the subspace defined by  $U_k$  contains those  $r + 1$  vectors  $z_{j^*}, z'_{j^*}, z_{j_1}, \dots, z_{j_{r-1}}$ , it must have dimension at least  $r + 1$ , which leads to a contradiction. We show this by showing that those  $r + 1$  vectors are linearly independent. Notice that the last  $r$  entries of  $z_{j^*}, z'_{j^*}$  are identical. The first  $\bar{D}$  entries of them should be different. Let  $z_{j^*} = [x \ z_{j^*}^-]^\top$  and  $z'_{j^*} = [x' \ z_{j^*}^-]^\top$ , and we know  $x \neq x'$ .

Suppose  $z_{j^*}, z'_{j^*}, z_{j_1}, \dots, z_{j_{r-1}}$  are linearly dependent. There exists  $c_{j^*}, c'_{j^*}, c_1, \dots, c_{r-1}$  such that  $c_{j^*} z_{j^*} + c'_{j^*} z'_{j^*} + c_1 z_{j_1} + \dots + c_{r-1} z_{j_{r-1}} = 0$ , where  $c_{j^*}, c'_{j^*}, c_1, \dots, c_{r-1}$  are  $r + 1$  real numbers such that at least one of them is nonzero. Since this equality implies  $(c_{j^*} + c'_{j^*}) z_{j^*}^- + c_1 z_{j_1}^- + \dots + c_{r-1} z_{j_{r-1}}^- = 0$  and  $z_{j^*}^-, z_{j_1}^-, \dots, z_{j_{r-1}}^-$  are linearly independent, we know that  $c_{j^*} = -c'_{j^*}$  and  $c_1 = \dots = c_{r-1} = 0$ . However, since  $x \neq x'$ , the first  $\bar{D}$  entries of  $c_{j^*} z_{j^*} + c'_{j^*} z'_{j^*} + c_1 z_{j_1} + \dots + c_{r-1} z_{j_{r-1}}$  cannot be all 0's, which is a contradiction. This proves the first statement in the proposition.

The implication from the first statement to the second is straightforward, since intersecting  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  implies intersecting  $\mathcal{P}^{i_1, \dots, i_{r-1}}$ , and the intersection between  $U_k$  and  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  being a single point implies that the intersection between  $U_k$  and  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  can at most be a single point.  $\square$

The proposition below shows that, when  $S$  and  $T$  are sufficiently large, if a subspace  $U_k$  intersects at least  $M^{r-2} + 1$  hyper-balls, then the  $i$ -th column of  $U_k$  is very close to the

$(\bar{D} + i)$ -th unit vector, up to a change of orthonormal basis.

**Proposition 10.** *Suppose  $S > r(2d+1)^3 + 2rMd^2$  and  $T > S + 4rMd^2$ . If a subspace  $U \in \mathbb{R}^{D \times r}$  intersects at least  $M^{r-2} + 1$  hyper-balls from  $\{\mathcal{S}^{i_1, \dots, i_{r-1}} : i_1, \dots, i_{r-1} \in \{0, 1, \dots, M-1\}\}$ , then, up to a change of orthonormal basis, we have  $U = [u_1 \ \dots \ u_r]$  where*

- $u_1 = [\kappa_1 \ \tau \ 0 \ \dots \ 0]$  where  $\kappa_1 \in \mathbb{R}^{\bar{D}}$  satisfies  $\|\kappa_1\| < \frac{rMd}{T}$  and  $\tau > \sqrt{1 - (\frac{rMd}{T})^2}$ , and
- for each  $i = 2, \dots, r$ ,  $u_i = [\kappa_i \ \xi_{i0} \ \xi_{i1} \ \dots \ \xi_{i(i-2)} \ \sigma \ 0 \ \dots \ 0]$  where  $\kappa_i \in \mathbb{R}^{\bar{D}}$  satisfies  $\|\kappa_i\| < \frac{2d+1}{S}$ ,  $\xi_{i0} \in \mathbb{R}$  satisfies  $|\xi_{i0}| < \frac{4rMd^2}{ST}$ , each  $\xi_{ij} \in \mathbb{R}$  (with  $1 \leq j \leq i-2$ ) satisfies  $|\xi_{ij}| < \frac{(2d+1)^2}{S^2}$ , and  $\sigma > \sqrt{1 - \frac{(2d+1)^2}{S^2} - \frac{16r^2M^2d^4}{S^2T^2} - \frac{(i-2)(2d+1)^4}{S^4}}$ .

*Proof.* The high level ideas of this proof is as follows: firstly, we find  $r$  linearly independent vectors  $v_1, \dots, v_r$  that  $U$  must contain; secondly, we show that we can orthonormalise those  $r$  vectors to  $u_1, \dots, u_r$  satisfying the requirements in the proposition. Since  $U$  is an  $r$ -dimensional subspace, and  $UU^\top$  is invariant up to a change of orthonormal basis, these will implies our proposition.

To construct  $v_1, \dots, v_r$ , we first show that for each  $j = 1, \dots, r-1$ , there exist  $r$  positive integers  $i_1^{(j)}, \dots, i_{j-1}^{(j)}, i_j, i'_j, i_{j+1}^{(j)}, \dots, i_{r-1}^{(j)} \in \{0, 1, \dots, M-1\}$  with  $i_j \neq i'_j$  such that  $U$  intersects both  $\mathcal{S}^{i_1^{(j)}, \dots, i_{j-1}^{(j)}, i_j, i_{j+1}^{(j)}, \dots, i_{r-1}^{(j)}}$  and  $\mathcal{S}^{i_1^{(j)}, \dots, i_{j-1}^{(j)}, i'_j, i_{j+1}^{(j)}, \dots, i_{r-1}^{(j)}}$ . Intuitively, this claim says that you can find at least two hyper-balls along each of the  $r-1$  directions that  $U$  intersects. This is easy to show straightforwardly. Suppose for certain  $j$  we cannot find two balls intersected by  $U$  that only differ in the  $j$ -th ‘‘coordinate’’  $i_j$ . Then all the intersected balls must have different values for the remaining  $r-2$  ‘‘coordinates’’:  $i_1, \dots, i_{j-1}, i_{j+1}, \dots, i_{r-1}$ . However, there can be at most  $M^{r-2}$  such balls, which contradicts to that  $U$  intersects at least  $M^{r-2} + 1$  balls.

With this claim, we are ready to construct  $v_1, \dots, v_r$ . For each  $j = 2, \dots, r$ , define  $v_j$  as follows. Find the two balls  $\mathcal{S}^{i_1^{(j-1)}, \dots, i_{j-2}^{(j-1)}, i_{j-1}, i_j^{(j-1)}, \dots, i_{r-1}^{(j-1)}}$  and  $\mathcal{S}^{i_1^{(j-1)}, \dots, i_{j-2}^{(j-1)}, i'_{j-1}, i_j^{(j-1)}, \dots, i_{r-1}^{(j-1)}}$  that  $U$  intersects. Then the two intersection points (notice that Proposition 9 makes sure each intersection is a point) must be

$$\begin{bmatrix} \bar{x}_{j1} & T & i_1^{(j-1)}S & \dots & i_{j-2}^{(j-1)}S & i_{j-1}S & i_j^{(j-1)}S & \dots & i_{r-1}^{(j-1)}S \end{bmatrix}$$

and

$$\begin{bmatrix} \bar{x}_{j2} & T & i_1^{(j-1)}S & \dots & i_{j-2}^{(j-1)}S & i'_{j-1}S & i_j^{(j-1)}S & \dots & i_{r-1}^{(j-1)}S \end{bmatrix}$$

for certain  $\bar{x}_{j1}, \bar{x}_{j2} \in \mathbb{R}^{\bar{D}}$  with two norms bounded by  $d$ . Let  $v_j$  be the difference between this two intersection points divided by  $(i_{j-1} - i'_{j-1})$  (which should also be contained in  $U$ ):

$$v_j = [\bar{x}_j \ 0 \ 0 \ \dots \ 0 \ S \ 0 \ \dots \ 0],$$

where  $\bar{x}_j := \frac{1}{i_{j-1} - i'_{j-1}}(\bar{x}_{j1} - \bar{x}_{j2}) \in \mathbb{R}^{\bar{D}}$  satisfies  $\|\bar{x}_j\| \leq \frac{2d}{i_{j-1} - i'_{j-1}} \leq 2d$  by the triangle inequality. To construct  $v_1$ , we take an arbitrary ball  $\mathcal{S}^{i_1^\dagger, \dots, i_{r-1}^\dagger}$  that  $U$  intersect. Let  $v_1'$  be the point of intersection (or the vector pointing from the origin to the point of intersection). We have

$$v_1' = \begin{bmatrix} \bar{x}^\dagger & T & i_1^\dagger \cdot S & \cdots & i_{r-1}^\dagger \cdot S \end{bmatrix}$$

for certain  $x^\dagger \in \mathbb{R}^{\bar{D}}$  such that  $\|\bar{x}^\dagger\| \leq d$ . Let

$$v_1 = v_1' - \sum_{j=2}^r i_{j-1}^\dagger \cdot v_j = [\bar{x}_1 \quad T \quad 0 \quad \cdots \quad 0],$$

where  $\bar{x}_1 = \bar{x}^\dagger - \sum_{j=2}^r i_{j-1}^\dagger \cdot \bar{x}_j$ . We have  $\|\bar{x}_1\| \leq d + \sum_{j=2}^r i_{j-1}^\dagger \cdot d < rMd$ . Moreover, since  $v_1', v_2, \dots, v_r$  belong to the subspace  $U$ ,  $v_1$ , being a linear combination of those, also belongs to  $U$ .

Next, we transform  $v_1, \dots, v_r$  to orthonormal  $u_1, \dots, u_r$ . Intuitively, when  $S$  and  $T$  are very large,  $v_1, \dots, v_r$  are almost orthogonal such that  $v_1 \approx Te_{\bar{D}+1}$  and  $v_i \approx Se_{\bar{D}+i}$  for  $i \geq 2$ . After the Gram-Schmidt process, those  $r$  vectors are only changed by a little bit, and the  $(\bar{D}+i)$ -th coordinate of  $u_i$  is almost 1, which implies the proposition. The detailed calculations are as follows.

We will show that, after orthogonalization but before normalization,  $u_1, \dots, u_r$  are as follows:

$$\begin{aligned} u_1 &= [x_1 \quad T \quad 0 \quad 0 \quad 0 \quad \cdots \quad 0] \\ u_2 &= [x_2 \quad \xi_{20} \quad S \quad 0 \quad 0 \quad \cdots \quad 0] \\ u_3 &= [x_3 \quad \xi_{30} \quad \xi_{31} \quad S \quad 0 \quad \cdots \quad 0] \\ u_4 &= [x_4 \quad \xi_{40} \quad \xi_{41} \quad \xi_{42} \quad S \quad \cdots \quad 0] \\ &\quad \vdots \\ u_r &= [x_r \quad \xi_{r0} \quad \xi_{r1} \quad \xi_{r2} \quad \xi_{r3} \quad \cdots \quad S] \end{aligned}$$

where

- $x_1 \in \mathbb{R}^{\bar{D}}$  satisfies  $\|x_1\| < rMd$ ,
- $x_i \in \mathbb{R}^{\bar{D}}$  for each  $i \geq 2$  satisfies  $\|x_i\| < 2d + 1$ ,
- $|\xi_{i0}| < \frac{4rMd^2}{T}$  for each  $i = 2, \dots, r$ , and
- $|\xi_{ij}| < \frac{(2d+1)^2}{S}$  for each  $i = 2, \dots, r$  and each  $j$  with  $1 \leq j \leq i - 2$ .

This can be shown by induction. For the base step, we know  $u_1 = v_1$ , so  $\|x_1\| = \|\bar{x}_1\| < rMd$ . For  $u_2$ , since  $u_2 = v_2 - \frac{v_2^\top u_1}{\|u_1\|^2} u_1$  in the Gram-Schmidt process, we have

$$\|x_2\| = \left\| \bar{x}_2 - \frac{\bar{x}_2^\top x_1}{\|x_1\|^2 + T^2} x_1 \right\| < \|\bar{x}_2\| + \frac{\|\bar{x}_2\| \|x_1\|^2}{T^2} \leq 2d + \frac{2d \times rMd}{T^2} < 2d + 1,$$

and

$$|\xi_{20}| = \left| 0 - \frac{\bar{x}_2^\top x_1}{\|x_1\|^2 + T^2} T \right| < \frac{\|\bar{x}_2\| \|x_1\|}{T} \leq \frac{2d \times rMd}{T} < \frac{4rMd^2}{T}.$$

Since the remaining coordinates of  $u_1$  are 0, the remaining coordinates of  $u_2$  are identical to those of  $v_2$ . This concludes the base step.

For the inductive step, suppose our claim holds for  $u_1, \dots, u_i$ . We consider  $u_{i+1}$ . Firstly, we only need to check for  $x_{i+1}$  and  $\xi_{(i+1)0}, \dots, \xi_{(i+1)(i-1)}$ , as the remaining coordinates of  $u_{i+1}$  are identical to those in  $v_{i+1}$  since these coordinates for each of  $u_1, \dots, u_i$  are 0 by the induction hypothesis. For  $x_{i+1}$ , we have

$$\begin{aligned} \|x_{i+1}\| &= \left\| \bar{x}_{i+1} - \sum_{s=1}^i \frac{v_{i+1}^\top u_s}{\|u_s\|^2} x_s \right\| \\ &= \left\| \bar{x}_{i+1} - \sum_{s=1}^i \frac{\bar{x}_{i+1}^\top x_s}{\|x_s\|^2 + \xi_{s0}^2 + \dots + \xi_{s(s-2)}^2 + S^2} x_s \right\| \\ &< 2d + \sum_{s=1}^i \frac{2d\|x_s\|^2}{S^2} && \text{(since } \|\bar{x}_{i+1}\| \leq 2d) \\ &< 2d + \frac{2d \cdot rMd}{S^2} + \frac{2d}{S^2} \cdot (i-1)(2d+1)^2 && \text{(induction hypothesis)} \\ &< 2d + 1. && \text{(since } S^2 > S > r(2d+1)^3 + 2rMd^2) \end{aligned}$$

For  $\xi_{(i+1)0}$ , we have

$$\begin{aligned} \|\xi_{(i+1)0}\| &= \left\| 0 - \frac{v_{i+1}^\top u_1}{\|u_1\|^2} T - \sum_{s=2}^i \frac{v_{i+1}^\top u_s}{\|u_s\|^2} \xi_{s0} \right\| \\ &= \left\| -\frac{\bar{x}_{i+1}^\top x_1}{\|x_1\|^2 + T^2} T - \sum_{s=2}^i \frac{\bar{x}_{i+1}^\top x_s}{\|x_s\|^2 + \xi_{s0}^2 + \dots + \xi_{s(s-2)}^2 + S^2} \xi_{s0} \right\| \\ &< \frac{2d\|x_1\|}{T} + \sum_{s=2}^i \frac{2d\|x_s\|}{S^2} |\xi_{s0}| && \text{(since } \|\bar{x}_{i+1}\| \leq 2d) \\ &< \frac{2rMd^2}{T} + \frac{r \cdot 2d(2d+1)}{S^2} \frac{4rMd^2}{T} && \text{(by induction hypothesis)} \\ &< \frac{2rMd^2}{T} + \frac{1}{S(2d+1)} \frac{4rMd^2}{T} && \text{(since } S > r(2d+1)^3) \\ &< \frac{2rMd^2}{T} + \frac{2rMd^2}{T} && \text{(obviously } S(2d+1) > 2) \\ &= \frac{4rMd^2}{T}. \end{aligned}$$



For each  $\xi_{(i+1)j}$  with  $1 \leq j \leq i-1$ , we have

$$\begin{aligned}
& \|\xi_{(i+1)j}\| \\
&= \left\| 0 - \frac{v_{i+1}^\top u_{j+1}}{\|u_{j+1}\|^2} S - \sum_{s=j+2}^i \frac{v_{i+1}^\top u_s}{\|u_s\|^2} \xi_{sj} \right\| \quad (\text{the last term vanishes if } i < j+2) \\
&= \left\| -\frac{\bar{x}_{i+1}^\top x_{j+1} \cdot S}{\|x_{j+1}\|^2 + \xi_{(j+1)0}^2 + \cdots + \xi_{(j+1)(j-1)}^2 + S^2} - \sum_{s=j+2}^i \frac{\bar{x}_{i+1}^\top x_s \cdot \xi_{sj}}{\|x_s\|^2 + \xi_{s0}^2 + \cdots + \xi_{s(s-2)}^2 + S^2} \right\| \\
&< \frac{2d(2d+1)}{S} + \frac{r \cdot 2d(2d+1)}{S^2} \frac{(2d+1)^2}{S} \\
&< \frac{2d(2d+1)}{S} + \frac{1}{r(2d+1)^4} \frac{(2d+1)^2}{S} \quad (\text{since } S > r(2d+1)^3) \\
&< \frac{2d(2d+1)}{S} + \frac{1}{S} < \frac{(2d+1)^2}{S}.
\end{aligned}$$

Finally, by normalizing  $u_1, \dots, u_r$  to unit vectors, the proposition follows easily.  $\square$

Proposition 9 and Proposition 10 imply the following corollary.

**Corollary 11.** *Suppose  $S > r(2d+1)^3 + 2rMd^2$  and  $T > S + 4rMd^2$ . If a subspace  $U \in \mathbb{R}^{D \times r}$  intersects at least  $M^{r-2} + 1$  hyper-balls from  $\{\mathcal{S}^{i_1, \dots, i_{r-1}} : i_1, \dots, i_{r-1} \in \{0, 1, \dots, M-1\}\}$ , then  $U$  intersects each hyperplane  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  at exactly one point.*

*Proof.* Proposition 10 shows that  $U = [u_1 \ \cdots \ u_r]$ , where  $u_1, \dots, u_r$  are as defined in the statement of Proposition 10. A point in  $U$  can be represented as  $x = c_1 u_1 + \cdots + c_r u_r$  for certain  $c_1, \dots, c_r \in \mathbb{R}$ . It is also in  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  if the last  $r$  coordinates of  $x$  are  $T, i_1 S, \dots, i_{r-1} S$  respectively. Therefore, the intersection between  $U$  and  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  are exactly those  $x = c_1 u_1 + \cdots + c_r u_r$  subjecting to

$$c_1 \begin{bmatrix} \tau \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} + c_2 \begin{bmatrix} \xi_{20} \\ \sigma \\ 0 \\ \vdots \\ 0 \end{bmatrix} + c_3 \begin{bmatrix} \xi_{30} \\ \xi_{31} \\ \sigma \\ \vdots \\ 0 \end{bmatrix} + \cdots + c_r \begin{bmatrix} \xi_{r0} \\ \xi_{r1} \\ \xi_{r2} \\ \vdots \\ \sigma \end{bmatrix} = \begin{bmatrix} T \\ i_1 S \\ i_2 S \\ \vdots \\ i_{r-1} S \end{bmatrix}.$$

There exists a unique  $(c_1, \dots, c_r)$  satisfying above equation, as the matrix

$$\begin{bmatrix} \tau & \xi_{20} & \xi_{30} & \cdots & \xi_{r0} \\ 0 & \sigma & \xi_{31} & \cdots & \xi_{r1} \\ 0 & 0 & \sigma & \cdots & \xi_{r2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & \sigma \end{bmatrix},$$

having determinant  $\tau\sigma^{r-1} > 0$ , is invertible. This proves that the intersection between  $U$  and  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  exists and is unique.  $\square$

Next, we show that each  $U_k$  must intersect at least  $M^{r-2} + 1$  hyper-balls in an optimal solution.

**Proposition 12.** *Let  $\delta = \min_{i,j \in \{1, \dots, \bar{N}\}} \frac{1}{2} \|\bar{x}_i - \bar{x}_j\|$  be the half of the minimum distance between two data points in the  $K$ -means instance. Suppose  $M > 4\bar{N}^2 d^2 / \delta^2$ , and  $M, S, T$  have sizes that is polynomial in  $N, K$  and  $D$ . If there exists  $U_k$  who intersect at most  $M^{r-2}$  hyper-balls in a solution  $(\mathcal{C}, \mathcal{U})$ , then there is a polynomial time algorithm that finds a replacement of  $U_k$ , the subspace  $U'_k \in \mathbb{R}^{D \times r}$ , such that the KSS objective is strictly increased, where  $U'_k$  intersects all the hyper-balls.*

*Proof.* Recall that for each datum  $\bar{x}_n$  in the  $K$ -means instance, we have constructed  $M^{r-1}$  data  $X_n = \{x_n^{i_1, \dots, i_{r-1}}\}$ . Given the solution  $\mathcal{U} = \{U_1, \dots, U_K\}$  we are considering, we say a datum  $x_n^{i_1, \dots, i_{r-1}}$  is *covered* if there exists  $U_k$  such that the distance between  $x_n^{i_1, \dots, i_{r-1}}$  and  $U_k$  is less than  $\delta$ :  $\|x_n^{i_1, \dots, i_{r-1}} - U_k U_k^\top x_n^{i_1, \dots, i_{r-1}}\| < \delta$ . As a result, if a datum is uncovered, it contributes at least  $\delta^2$  to the KSS objective.

Firstly, we show that there exists  $n \in \{1, \dots, \bar{N}\}$  such that  $X_n$  contains at least  $\frac{\bar{N}-K}{\bar{N}} M^{r-1}$  uncovered data<sup>1</sup>. To see this,  $\{X_n : n = 1, \dots, \bar{N}\}$  form an equal-sized partition of all the  $N = M^{r-1} \bar{N}$  data. By Proposition 9 and our definition of  $\delta$ , each subspace that intersects at least  $M^{r-2} + 1$  balls can at most cover a single datum in each ball. Let  $K_1$  be the number of subspaces that intersect at least  $M^{r-2} + 1$  balls, and let  $K_2 = K - K_1$  be the number of subspaces that intersect at most  $M^{r-2}$  balls. The total number of covered data is at most

$$K_1 M^{r-1} \cdot 1 + K_2 M^{r-2} \cdot \bar{N} \leq K M^{r-1}.$$

By the pigeonhole principle, there exists  $X_n$  that contains at most  $\frac{K M^{r-1}}{\bar{N}}$  covered data, which implies our claim. Let  $X_{n^*}$  be such a partition.

Next, we show that the subspace that contains all the  $M^{r-1}$  data in  $X_{n^*}$ , which is

$$U'_k = \begin{bmatrix} \frac{\bar{x}_{n^*}}{\sqrt{\|\bar{x}_{n^*}\|^2 + T^2}} & \mathbf{0} & \mathbf{0} & \cdots & \mathbf{0} \\ \frac{T}{\sqrt{\|\bar{x}_{n^*}\|^2 + T^2}} & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \end{bmatrix},$$

satisfies the requirement in the proposition (it is easy to see that  $U'_k$  is orthonormal and  $U'_k (U'_k)^\top x_{n^*}^{i_1, \dots, i_{r-1}} = x_{n^*}^{i_1, \dots, i_{r-1}}$ ). When replacing  $U_k$  by  $U'_k$  for certain  $U_k$  that intersects at

<sup>1</sup>It is reasonable to assume  $\bar{N} > \bar{K} = K$  in the  $K$ -means instance, for otherwise the  $K$ -means instance is trivial: just let all the data be the centers.

most  $M^{r-2}$  balls, the loss in the objective function is at most  $M^{r-2}\bar{N} \cdot 4d^2$  (even assuming  $U_k$  contains all the  $M^{r-2}\bar{N}$  data in those  $M^{r-2}$  balls it intersects, these data can be assigned to  $U'_k$  with cost at most  $(2d)^2$  as  $U'_k$  now intersects all the hyper-balls), and the increment in the objective function is at least  $\frac{\bar{N}-K}{N}M^{r-1} \cdot \delta^2 \geq \frac{1}{N}M^{r-1}\delta^2$  (those  $\frac{\bar{N}-K}{N}M^{r-1}$  data not covered, each of which contributing at least  $\delta^2$  to the KSS objective, will have zero contribution if  $U'_k$  replaces  $U_k$ ). By our assumption  $M > 4\bar{N}^2d^2/\delta^2$ , the benefit of replacing  $U_k$  by  $U'_k$  is more than the loss.

Finally, to conclude the proof, it is obvious that finding such  $U'_k$  requires polynomial time.  $\square$

The proposition below is an extension of Lemma 6 to higher dimension, which says that the distance between  $x \in \mathcal{S}^{i_1, \dots, i_{r-1}}$  and the subspace  $U$  is close to the distance between  $x$  and the intersection  $s$  of  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  and  $U$ . The former distance corresponds to the KSS objective, while the latter distance corresponds to the  $K$ -means objective if viewing all data in  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  as a  $K$ -means instance and  $s$  as a center.

**Proposition 13.** *Let  $U \in \mathbb{R}^{D \times r}$  be a subspace that intersects at least  $M^{r-2} + 1$  hyper-balls. Fix an arbitrary hyperplane  $\mathcal{P} := \mathcal{P}^{i_1, \dots, i_{r-1}}$ , and let  $s \in \mathbb{R}^D$  be the unique intersection between  $\mathcal{P}$  and  $U$  (Corollary 11). For any  $\zeta > 0$ , there exists  $\Pi > 0$  that depends only on  $\zeta, r$  and  $d$  such that for any  $x \in \mathcal{S}^{i_1, \dots, i_{r-1}}$ , any  $M \geq \Pi$ ,  $S = M(r(2d+1)^3 + 2rMd^2)$ , and  $T = MS$ , we have*

$$1 \leq \frac{\|x - s\|}{\|x - UU^\top x\|} \leq 1 + \zeta.$$

*Proof.* Proposition 10 shows that  $U = [u_1 \ \dots \ u_r]$ , where  $u_1, \dots, u_r$  are as defined in the statement of Proposition 10. The left part of the inequality is trivial, as the point in  $U$  that is closest to  $x$  is precisely  $x$ 's projection on  $U$ , namely,  $UU^\top x$ . Thus, it remains to show the right part. By the definition of limit, this is equivalent as

$$\lim \frac{\|x - s\|}{\|x - UU^\top x\|} = 1 \quad \text{or} \quad \lim \|x - s\| = \lim \|x - UU^\top x\|,$$

where all the limits above, as well as all the limits in this proof, are taken over  $M \rightarrow \infty$  while fixing  $S = M(r(2d+1)^3 + 2rMd^2)$  and  $T = MS$ . It is important to notice that the limiting behavior should not depend on  $\bar{D}$ , as  $\Pi$  can only depend on  $\zeta, r$  and  $d$  (it can depend on  $D$ , as  $D = \bar{D} + r$  is dependent on  $r$ ). We will make sure this is true, by ensuring our evaluation of the limit based solely on the upper bounds of those  $\|\kappa_i\|$ 's and  $|\xi_{ij}|$ 's given in Proposition 10. In particular, although the dimension of  $\kappa_i$ 's depends on (in fact, is exactly)  $\bar{D}$ , but the bounds of their 2-norms, which is  $\frac{rMd}{T}$  for  $\kappa_1$  and  $\frac{2d+1}{S}$  for remaining  $\kappa_i$ , do not depend on  $\bar{D}$ .

Since  $|\|x - s\| - \|x - UU^\top x\|| \leq \|s - UU^\top x\|$ , it suffices to show that  $\lim \|s - UU^\top x\| = 0$ . Let  $B = [\kappa_1 \ \dots \ \kappa_r] \in \mathbb{R}^{\bar{D} \times r}$ ,  $\bar{x} \in \mathbb{R}^{\bar{D}}$  be the first  $\bar{D}$  coordinates of

$x, \underline{x} = [T \ i_1 S \ \cdots \ i_{r-1} S]$  be the last  $r$  coordinate of  $x$ , and

$$A = \begin{bmatrix} \tau & \xi_{20} & \xi_{30} & \cdots & \xi_{r0} \\ 0 & \sigma & \xi_{31} & \cdots & \xi_{r1} \\ 0 & 0 & \sigma & \cdots & \xi_{r2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & \sigma \end{bmatrix},$$

so we have

$$U = \begin{bmatrix} B \\ A \end{bmatrix} \quad \text{and} \quad x = \begin{bmatrix} \bar{x} \\ \underline{x} \end{bmatrix}.$$

We express  $UU^\top x$  and  $s$  in terms of  $A, B, \bar{x}$  and  $\underline{x}$ . It is easy to compute  $UU^\top x$ :

$$UU^\top x = \begin{bmatrix} BB^\top & BA^\top \\ AB^\top & AA^\top \end{bmatrix} \begin{bmatrix} \bar{x} \\ \underline{x} \end{bmatrix} = \begin{bmatrix} BB^\top \bar{x} + BA^\top \underline{x} \\ AB^\top \bar{x} + AA^\top \underline{x} \end{bmatrix}.$$

For  $s$ , recall that  $s = \sum_{j=1}^r c_j u_j$  where  $c_1, \dots, c_j$  satisfies

$$c_1 \begin{bmatrix} \tau \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} + c_2 \begin{bmatrix} \xi_{20} \\ \sigma \\ 0 \\ \vdots \\ 0 \end{bmatrix} + c_3 \begin{bmatrix} \xi_{30} \\ \xi_{31} \\ \sigma \\ \vdots \\ 0 \end{bmatrix} + \cdots + c_r \begin{bmatrix} \xi_{r0} \\ \xi_{r1} \\ \xi_{r2} \\ \vdots \\ \sigma \end{bmatrix} = \begin{bmatrix} T \\ i_1 S \\ i_2 S \\ \vdots \\ i_{r-1} S \end{bmatrix},$$

which is

$$\begin{bmatrix} c_1 \\ \vdots \\ c_r \end{bmatrix} = A^{-1} \underline{x}.$$

Therefore, we have

$$s = U \begin{bmatrix} c_1 \\ \vdots \\ c_r \end{bmatrix} = \begin{bmatrix} B \\ A \end{bmatrix} A^{-1} \underline{x} = \begin{bmatrix} BA^{-1} \underline{x} \\ AA^{-1} \underline{x} \end{bmatrix}.$$

Putting together,

$$UU^\top x - s = \begin{bmatrix} BB^\top \bar{x} + BA^\top \underline{x} \\ AB^\top \bar{x} + AA^\top \underline{x} \end{bmatrix} - \begin{bmatrix} BA^{-1} \underline{x} \\ AA^{-1} \underline{x} \end{bmatrix} = UB^\top \bar{x} + U(A^\top - A^{-1}) \underline{x}.$$

To conclude the proposition, we will show

$$\lim \|UB^\top \bar{x}\| = 0 \quad \text{and} \quad \lim \|U(A^\top - A^{-1}) \underline{x}\| = 0.$$

It is easy to see the first limit by the Cauchy-Schwarz inequality:

$$\left\| UB^\top \bar{x} \right\| \leq \|U\|_F \|B^\top\|_F \|\bar{x}\|,$$

where  $\|\cdot\|_F$  denotes the Frobenius norm. We have  $\|U\|_F = \sqrt{r}$  (as the sum of the squares for each column of  $U$  is 1) and  $\|\bar{x}\| \leq d$  (by our definition of  $d$ ), which are both bounded. Moreover, by Proposition 10,

$$\lim \left\| B^\top \right\|_F^2 = \lim \left( \sum_{j=1}^r \|\kappa_j\|^2 \right) \leq \lim \left( \left( \frac{rMd}{T} \right)^2 + (r-1) \left( \frac{2d+1}{S} \right)^2 \right) = 0.$$

Therefore, we conclude the first limit.

It is a little bit more tricky to show the second limit, as  $\|\underline{x}\|$  tends to infinity. By Cauchy-Schwarz inequality again, we have

$$\left\| U \left( A^\top - A^{-1} \right) \underline{x} \right\| \leq \|U\|_F \left\| \left( A^\top - A^{-1} \right) \underline{x} \right\|.$$

Since  $\|U\|_F = \sqrt{r}$  is bounded, it remains to show that

$$\lim \left\| \left( A^\top - A^{-1} \right) \underline{x} \right\| = 0.$$

Let  $a_{ij}$  be the entry  $(i, j)$  of the matrix  $A^\top - A^{-1}$ . Noticing that  $A^\top$  is lower-triangular and  $A^{-1}$  is upper-triangular, we have  $a_{11} = \tau - \frac{1}{\tau}$ ,  $a_{22} = \dots = a_{rr} = \sigma - \frac{1}{\sigma}$ ,  $a_{ij} = \xi_{i(j-1)}$  for all  $i > j$ , and  $a_{ij}$  is the entry  $(i, j)$  of  $A^{-1}$  for all  $i < j$ . We will show the limit by the following two steps.

1. We first show that  $a_{i1} = o(\frac{1}{T})$  for all  $i = 1, \dots, r$ , and  $a_{ij} = o(\frac{1}{S})$  for all  $i \neq j$  and  $j \neq 1$ . Since the first entry of  $\underline{x}$  is  $T$ , and the remaining entries of  $\underline{x}$  is  $\Theta(S)$ , noticing that  $r$  is a constant, this will imply

$$\lim \left\| \left( A^\top - A^{-1} \right) \underline{x} \right\|^2 = \lim \left( (a_{11}T)^2 + \sum_{j=2}^r (a_{jj}i_{j-1}S)^2 \right).$$

2. In the second step, we show that the limit above is 0.

It is straightforward to show 2. Since Proposition 10 implies that  $0 < 1 - \tau^2 < (\frac{rMd}{T})^2$  and  $0 < 1 - \sigma^2 < (\frac{2d+1}{S})^2 + o(\frac{1}{S^2})$ , we have

$$\lim |a_{11}T| = \lim \frac{(1 - \tau^2)T}{\tau} \leq \frac{\lim (\frac{rMd}{T})^2 T}{1} = r^2 d^2 \lim \frac{M^2}{T} = 0,$$

as  $S = M(r(2d+1)^3 + 2rMd^2) > M^2 = \omega(M)$  and  $T = MS = \omega(M^2)$ . Similarly, for each  $j = 2, \dots, r$ , we have

$$\lim |a_{jj}i_{j-1}S| = \lim \frac{(1 - \sigma^2)i_{j-1}S}{\sigma} \leq \lim \left( \frac{2d+1}{S} \right)^2 i_{j-1}S = 0.$$

Putting together by noting that  $r$  is constant concludes 2.

To show 1, by noticing  $S = \omega(M)$ , Proposition 10 implies  $a_{i1} = \xi_{i0} = \Theta(\frac{M}{ST}) = o(\frac{1}{T})$  and  $a_{ij} = \xi_{i(j-1)} = \Theta(\frac{1}{S^2}) = o(\frac{1}{S})$  for all  $i < j$ . It remains to show that  $a_{ij} = o(\frac{1}{S})$  for all  $i > j$ , which requires analysis on the upper-right part of the matrix  $A^{-1}$ .

Let  $A_d$  be the diagonal matrix whose diagonal entries are identical to  $A$ , and  $A_u$  be the strictly upper-triangular matrix whose upper-right part is identical to  $A$ , so that  $A = A_d + A_u$ . The calculation in [3] reveals that

$$A^{-1} = \left( \sum_{s=0}^r (-A_d^{-1}A_u)^s \right) A_d^{-1}.$$

Since  $\tau^{-1} = \Theta(1)$  and  $\sigma^{-1} = \Theta(1)$ , the limit of  $A_d^{-1}$  is an identity matrix. Since the dimension of  $A$  is constant and the summation above also has constantly many terms, to analyze the asymptotic behavior of upper-right entries of  $A^{-1}$ , we can take  $A^{-1} = \sum_{s=0}^r (-A_u)^s$ . We know that each entry of  $A_u$ , which is certain  $\xi_{ij}$ , is  $o(\frac{1}{S})$ . Moreover, since the matrix has a constant dimension, it is easy to see that, if each entry in  $A_u^s$  is  $o(\varsigma)$  for certain infinitesimal  $\varsigma$ , then each entry in  $A_u^{s+1}$  is  $o(\varsigma^2)$ . As a result, we can conclude that each entry in the upper-right part of  $A^{-1}$  is  $o(\frac{1}{S})$ . Therefore, we conclude 2.  $\square$

Finally, we are ready to prove Theorem 8. We show that, by choosing an arbitrary  $\zeta$  with  $0 < \zeta < \sqrt{1+\epsilon} - 1$  and letting  $\varepsilon = \frac{1+\epsilon}{(1+\zeta)^2} - 1$ , a  $(1+\varepsilon)$ -approximation polynomial time algorithm for the KSS problem implies a  $(1+\varepsilon)$ -approximation polynomial time algorithm for the  $K$ -means problem.

Given a  $K$ -means instance, we construct the KSS instance based on our construction, with  $M = \max\{4\bar{N}^2d^2/\delta^2, \Pi\}$  where  $\Pi$  is given in Proposition 13,  $S = M(r(2d+1)^3 + 2rMd^2)$ , and  $T = MS$ . We first remark that the constructed KSS instance has a polynomial size. The reduction showing the APX-hardness of  $K$ -means in [1] satisfies  $d = \delta = \sqrt{2}$ .<sup>2</sup> Since  $1/\zeta$  depends only on  $\epsilon$  which is a constant,  $\Pi$ , depending only on  $1/\zeta, r$  and  $d$ , is also a constant, so  $M$  is of polynomial size. Since  $r$  is a constant,  $S$  and  $T$  are of polynomial size as well. The total number of data,  $N = M^{r-1}\bar{N}$ , is also of polynomial size.

Suppose there is a  $(1+\varepsilon)$ -approximation polynomial algorithm for the KSS problem. We can assume without loss of generality that each  $U_k$  intersects at least  $M^{r-2} + 1$  hyperballs: if the algorithm output certain  $U_k$  fail to satisfy this, we can apply Proposition 12 to

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<sup>2</sup>For a quick glance of the reduction in [1], the authors first show that the Vertex Cover problem on triangle-free graphs is APX-hard, then a datum  $x_{ij} = e_i + e_j$  is constructed for each edge  $(i, j)$  with  $1 \leq i, j \leq n$ . Therefore, each data point only have two non-zero entries which equal to 1.

find a better replacement of each such  $U_k$ , which only takes extra polynomial time. By our construction, the  $K$ -means instance is duplicated to  $M^{r-1}$  copies such that each  $\mathcal{S}^{i_1, \dots, i_{r-1}}$  contains a copy, and the solution  $\mathcal{U} = \{U_1, \dots, U_K\}$  induces a  $K$ -means solution to each of these copies. In particular, for each  $\mathcal{S}^{i_1, \dots, i_{r-1}}$ , each center  $s_k$  in the  $K$ -means problem is exactly the intersection of  $U_k$  and  $\mathcal{P}^{i_1, \dots, i_{r-1}}$  (Corollary 11 makes sure the intersection exists and is unique).

Given  $\mathcal{U}$ , let  $\phi_{KSS}(\mathcal{U})$  be the KSS objective value for solution  $(\mathcal{C}, \mathcal{U})$ , where  $C_k \in \mathcal{C}$  consists of all data which are closest to  $U_k$ . Let  $\phi_{KSS}^{i_1, \dots, i_{r-1}}(\mathcal{U})$  be the contributions to the KSS objective value from all data in  $\mathcal{S}^{i_1, \dots, i_{r-1}}$ , so that  $\phi_{KSS}(\mathcal{U}) = \sum_{i_1, \dots, i_{r-1}} \phi_{KSS}^{i_1, \dots, i_{r-1}}(\mathcal{U})$ . Let  $\phi_{KM}^{i_1, \dots, i_{r-1}}(\mathcal{U})$  be the  $K$ -means objective value for the  $K$ -means solution induced in the hyper-ball  $\mathcal{S}^{i_1, \dots, i_{r-1}}$ . Let  $\{s_1^*, \dots, s_K^*\}$  (with  $s_k^* \in \mathbb{R}^{\bar{D}}$ ) be an optimal solution to the  $K$ -means instance, and  $OPT_{KM}$  be the optimal objective value. Let  $\mathcal{U}^* = \{U_1^*, \dots, U_K^*\}$  where

$$U_k^* = \begin{bmatrix} \frac{s_k^*}{\sqrt{\|s_k^*\|^2 + T^2}} & \mathbf{0} & \mathbf{0} & \dots & \mathbf{0} \\ \frac{T}{\sqrt{\|s_k^*\|^2 + T^2}} & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 \end{bmatrix},$$

so  $\mathcal{U}^*$  induces the same optimal solution in each of the  $M^{r-1}$  duplicated  $K$ -means instances.

By Proposition 13, we have

$$\phi_{KSS}^{i_1, \dots, i_{r-1}}(\mathcal{U}^*) \leq \phi_{KM}^{i_1, \dots, i_{r-1}}(\mathcal{U}^*) \quad \text{and} \quad \phi_{KM}^{i_1, \dots, i_{r-1}}(\mathcal{U}) \leq (1 + \zeta)^2 \phi_{KSS}^{i_1, \dots, i_{r-1}}(\mathcal{U})$$

for each  $i_1, \dots, i_{r-1}$ , and we also have

$$\phi_{KSS}(\mathcal{U}) \leq (1 + \varepsilon) \phi_{KSS}(\mathcal{U}^{OPT}) \leq (1 + \varepsilon) \phi_{KSS}(\mathcal{U}^*)$$

where  $\mathcal{U}^{OPT}$  is an optimal solution to the constructed KSS instance. Putting together,

$$\begin{aligned} \sum_{i_1, \dots, i_{r-1}} \phi_{KM}^{i_1, \dots, i_{r-1}}(\mathcal{U}) &\leq \sum_{i_1, \dots, i_{r-1}} (1 + \zeta)^2 \phi_{KSS}^{i_1, \dots, i_{r-1}}(\mathcal{U}) \\ &= (1 + \zeta)^2 \phi_{KSS}(\mathcal{U}) \\ &\leq (1 + \zeta)^2 (1 + \varepsilon) \phi_{KSS}(\mathcal{U}^*) \\ &= (1 + \epsilon) \sum_{i_1, \dots, i_{r-1}} \phi_{KSS}^{i_1, \dots, i_{r-1}}(\mathcal{U}^*) \\ &\leq (1 + \epsilon) \sum_{i_1, \dots, i_{r-1}} \phi_{KM}^{i_1, \dots, i_{r-1}}(\mathcal{U}^*) \\ &= (1 + \epsilon) M^{r-1} OPT_{KM}. \end{aligned}$$

This shows that, the *average* of the  $K$ -means objective values induced in all the duplicated instances is at most  $(1 + \epsilon)OPT_{KM}$ . Therefore, we can exam the performances of the solutions induced in all those  $M^{r-1}$  identical  $K$ -means instance duplicates, and find the one with the best performance. This will gives a  $(1 + \epsilon)$ -approximation to the  $K$ -means instance.

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