

Lecture 4. Affine and convex sets

Affine combination: $\forall \theta_i \in \mathbb{R}, \theta_1 + \dots + \theta_n = 1, \underbrace{\theta_1 x_1 + \dots + \theta_n x_n}$

Affine set: given $x_1, \dots, x_n \in S$, \forall affine combination $\in S$.

Example. solution to a linear equation $S = \{x : Ax = b\}$.

$$(Ax_1 = Ax_2 = b \Rightarrow \forall \theta_1 + \theta_2 = 1, A(\theta_1 x_1 + \theta_2 x_2) = \theta_1 Ax_1 + \theta_2 Ax_2 = b)$$

Conversely, any affine set is the solution to a linear equation set.

Why? $S' = S - x_0$ is a linear space. $\forall x_1, x_2 \in S' \quad a_1 x_1 + a_2 x_2 \in S'$.

$$\frac{x_1 + x_0}{x_2 + x_0} \in S \Rightarrow a_1 x_1 + a_2 x_2 + x_0 = a_1(x_1 + x_0) + a_2(x_2 + x_0) + (1-a_1-a_2)x_0 \in S$$

In particular, if $A \in \mathbb{R}^{1 \times n}$. $S = \{x : w^T x = b\}$ is a hyperplane.

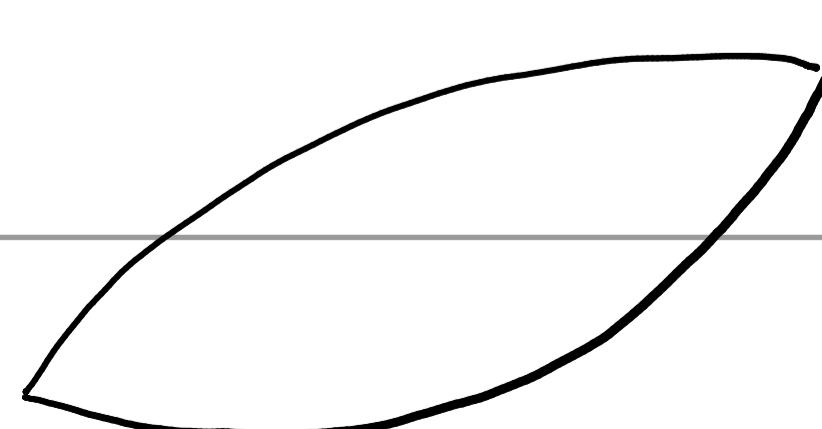
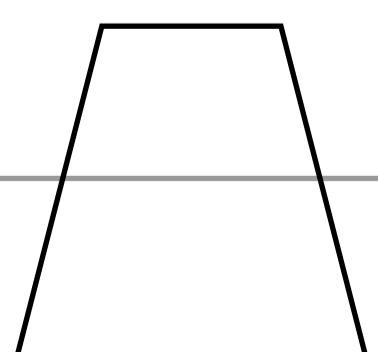
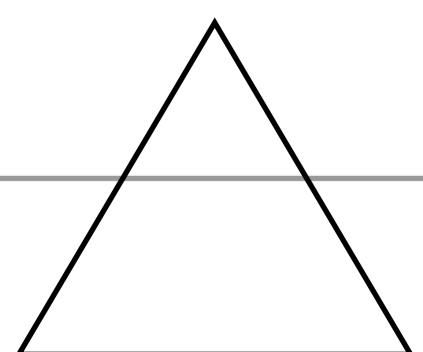
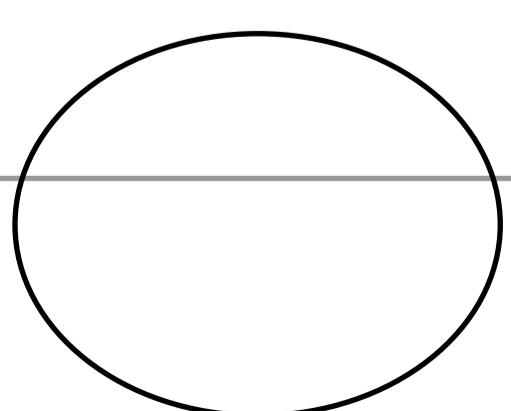
If $A \neq 0$, affine set $S \neq \mathbb{R}^n$ is the intersection of finite hyperplanes.

line: $z = x + \theta(y-x) = (1-\theta)x + \theta y, \theta \in \mathbb{R} \rightarrow$ affine sets

segment: $z = x + \theta(y-x) = (1-\theta)x + \theta y, \theta \in [0,1] \rightarrow$? convex sets.

Convex combination: $\theta_1 x_1 + \theta_2 x_2 + \dots + \theta_n x_n, \forall \theta_i \geq 0, \sum \theta_i = 1$.

Convex sets: If $x_1, \dots, x_n \in S$, any convex combination $\in S$, ($n=2$ suffices)

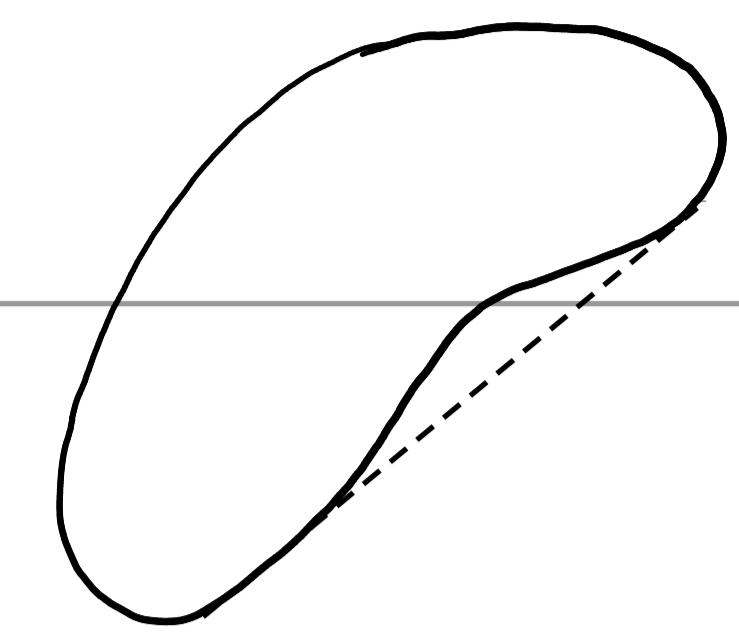


in particular

is not convex.

Convex hull: set of all convex combinations of points in S .

$$\text{conv}(S) \triangleq \left\{ \sum \theta_i x_i : \theta_i \geq 0, \sum \theta_i = 1, x_i \in S \right\}$$

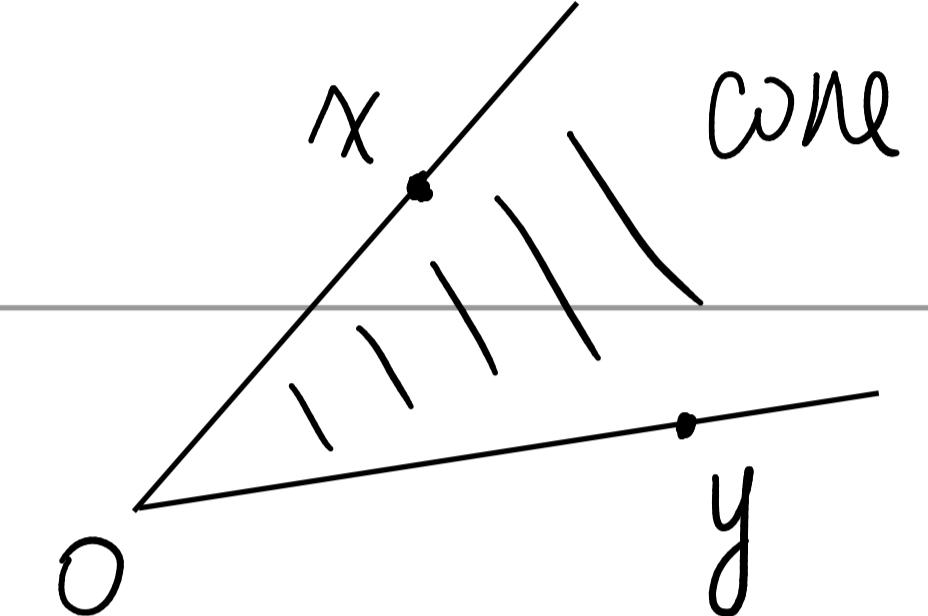


Conic combination: $z = \theta_1 x + \theta_2 y$. $\theta_1, \theta_2 \geq 0$.

Convex cone: set of all conic combinations of points in S .

Examples of convex sets:

\mathbb{R}^n affine sets hyperplanes

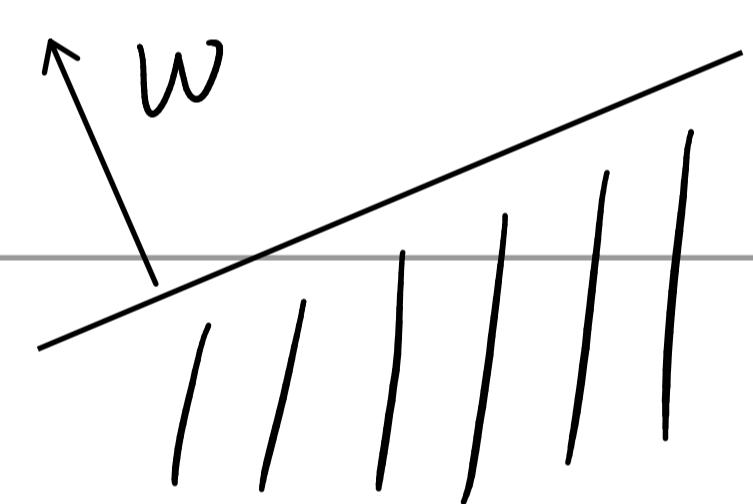


hyperplane $\{x : w^T x + b = 0\}$. halfspace $\{x : w^T x + b \leq 0\}$

Halfspaces are convex, but not affine. open halfspaces also convex

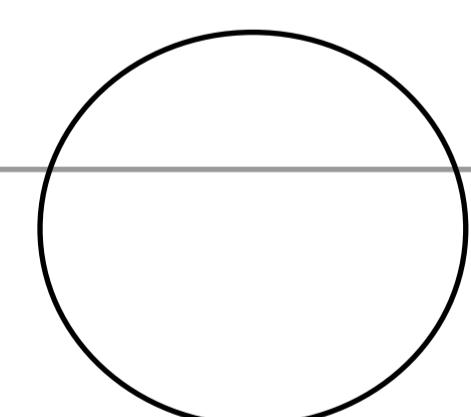
$\forall x, y \in \{x : w^T x + b \leq 0\}$. let $\bar{\theta} = 1 - \theta$

$$w^T(\theta x + \bar{\theta} y) + b = \theta(w^T x + b) + \bar{\theta}(w^T y + b) \leq 0$$



Euclidean balls and ellipsoids are convex.

Ball: $\{x : \|x - x_0\| \leq r\}$. $\forall x, y \in B(x_0, r)$.



$$\|\theta x + \bar{\theta} y - x_0\| = \|\theta(x - x_0) + \bar{\theta}(y - x_0)\| \leq \theta \|x - x_0\| + \bar{\theta} \|y - x_0\|$$

Ellipsoid: $\{(x_1, x_2)^T : \frac{x_1^2}{\lambda_1^2} + \frac{x_2^2}{\lambda_2^2} \leq 1\}$. convex. why? homework.

Proposition: The image of a convex set under an affine transform is

also convex. i.e., $C \subseteq \mathbb{R}^n$ convex $\Rightarrow AC + b = \{Ax + b : x \in C\}$ convex.

Proof. Let $f(x) = Ax + b$, $\forall y_1, y_2 \in f(C)$. $\exists x_i$, $y_i = f(x_i)$.

$$\theta y_1 + \bar{\theta} y_2 = \theta(Ax_1 + b) + \bar{\theta}(Ax_2 + b) = A(\theta x_1 + \bar{\theta} x_2) + b \in f(C) \quad \square$$

Question: How about inverse image? (homework) $\uparrow \in C$ by convexity.

Proposition: some other convexity-preserving operations.

- intersection: if C and D are convex, then so is $C \cap D$.
- Cartesian product: $C \times D \triangleq \{ (x_1, x_2) : x_1 \in C, x_2 \in D \}$ uncountably infinite many intersections ✓
- set sum (Minkowski addition): $C + D \triangleq \{ x_1 + x_2 : x_1 \in C, x_2 \in D \}$

A nongeometric example of convex sets: positive semidefinite matrices.

$$S_+^n \triangleq \{ A \in \mathbb{R}^{n \times n} : A \succeq 0 \} \quad S_{++}^n \triangleq \{ A \in \mathbb{R}^{n \times n} : A > 0 \} \text{ convex.}$$

Proof: $\forall A_1, A_2 \in S_+^n$. A_1, A_2 symmetric $\Rightarrow \theta A_1 + \bar{\theta} A_2$ symmetric

$$\forall v, v^T A_1 v \geq 0, v^T A_2 v \geq 0 \Rightarrow v^T (\theta A_1 + \bar{\theta} A_2) v \geq 0 \quad \square$$

Polyhedron (polyhedra) 多面体. polytope 多胞体.

A polyhedron is the intersection of some halfspaces.

$$\text{Halfspace: } \{ x : w^T x + b \leq 0 \}. \text{ Polyhedron: } \{ x : h_i. w_i^T x + b_i \leq 0 \}.$$

affine sets, halfspaces \subseteq polyhedra (may unbounded) \subseteq convex sets.

A polytope is an bounded polyhedron. Feasible sets of LP \subseteq polyhedra

Simplex (simplices / simplexes) 单纯形 : simplest polytope.

n -dimensional simplex : convex hull of ($n+1$) nondegenerate points.

Examples : 0-simplex : point 1-simplex : line segment.

2-simplex : triangle 3-simplex : tetrahedron.

nondegenerate ? $n+1$ affinely independent points in \mathbb{R}^n .

$$u_0 = \theta_1 u_1 + \cdots + \theta_n u_n \quad \sum \theta_i = 1 \iff \theta_1(u_1 - u_0) + \cdots + \theta_n(u_n - u_0) = 0.$$

u_0, u_1, \dots, u_n affinely dependent

$u_1 - u_0, \dots, u_n - u_0$ linearly dependent

$$(n+1)\text{-dimensional simplex } S = \{ \theta_0 u_0 + \cdots + \theta_n u_n : \sum \theta_i = 1, \theta_i \geq 0 \}$$

standard simplex. $u_0 = 0$, u_1, \dots, u_n = unit vector of n dimensions.

let $y \stackrel{\Delta}{=} (\theta_1, \dots, \theta_n)^T \in \mathbb{R}^n$. $B = (u_1 - u_0, \dots, u_n - u_0) \in \mathbb{R}^{n \times n}$ full rank.

$$S = \{ \theta_0 u_0 + \cdots + \theta_n u_n : \sum \theta_i = 1, \theta_i \geq 0 \} = \{ u_0 + B y : \sum y_i \leq 1, y_i \geq 0 \}$$

$$B \text{ has full rank} \Rightarrow \exists A = B^{-1} \in \mathbb{R}^n. \quad A(u_0 + B y) = A u_0 + y. \quad Hy.$$

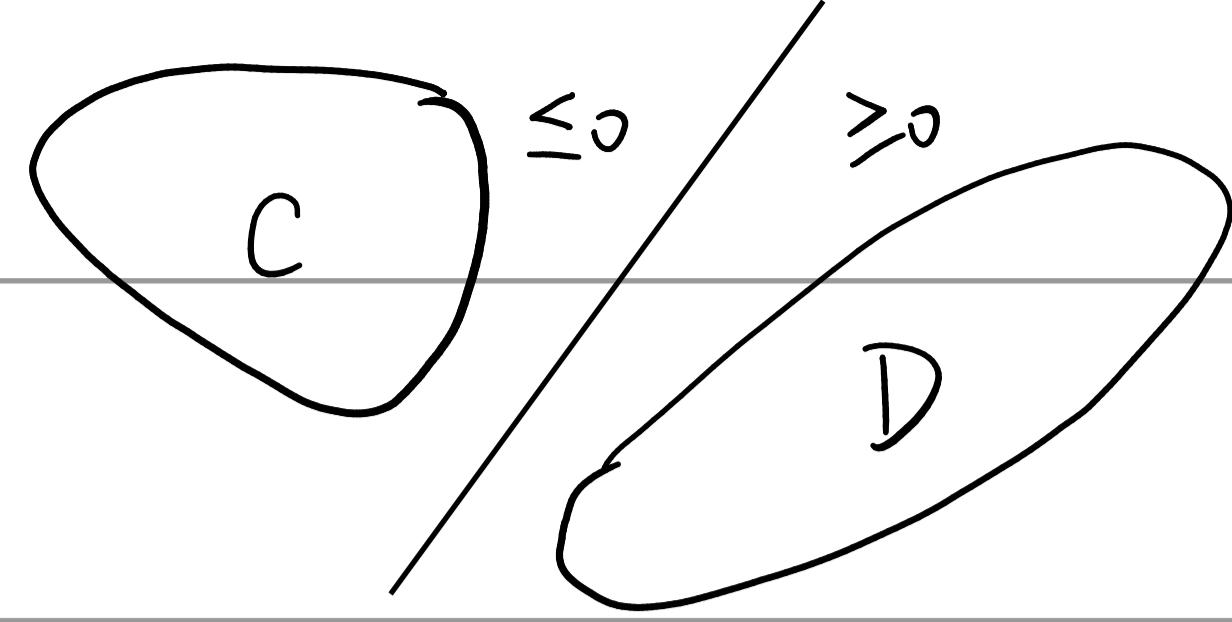
$$\text{So } \forall x \in \mathbb{R}^n, x \in S \text{ iff } \exists y, x = u_0 + B y \text{ iff } \exists y, Ax = Au_0 + y.$$

$$A = \begin{pmatrix} a_1^T \\ \vdots \\ a_n^T \end{pmatrix}. \quad y_i \geq 0 \Rightarrow a_i^T x \geq a_i^T u_0. \quad \sum y_i \leq 1 \Rightarrow \sum a_i^T (u_0 - x) \leq 1.$$

Separating hyperplane theorem and supporting hyperplane theorem.

(One of the most important and fundamental property of convex sets.)

Separating hyperplane theorem.



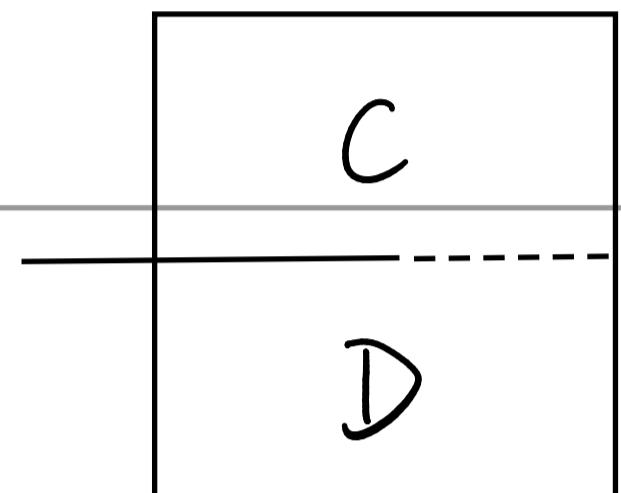
Let C, D be two disjoint convex sets.

Then \exists hyperplane $w^T x + b = 0$ separating C and D .
 $(w^T x + b \leq 0, x \in C)$
 $(w^T x + b \geq 0, x \in D)$

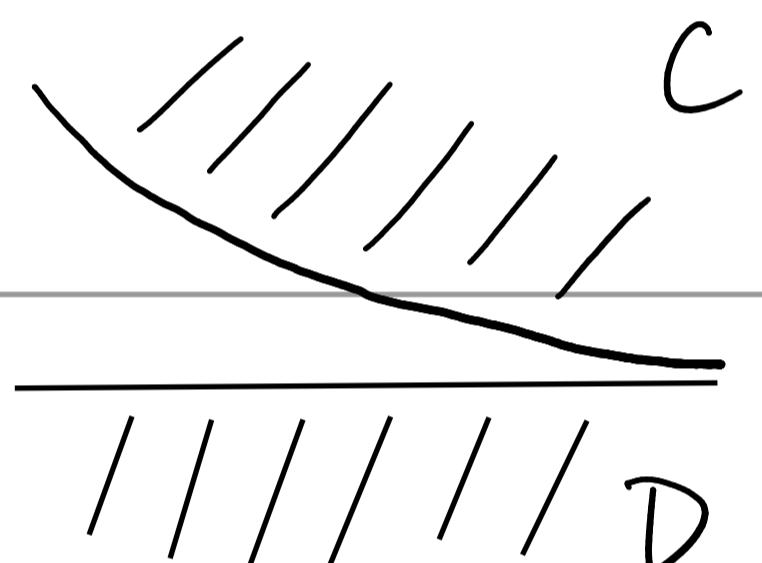
Strict separating theorem. (strict :
 $w^T x + b < 0, x \in C$
 $w^T x + b > 0, x \in D$)

Let C, D be two disjoint closed convex sets.

s.t. at least one of them is bounded.



Then \exists strict separating hyperplane $w^T x + b = 0$.

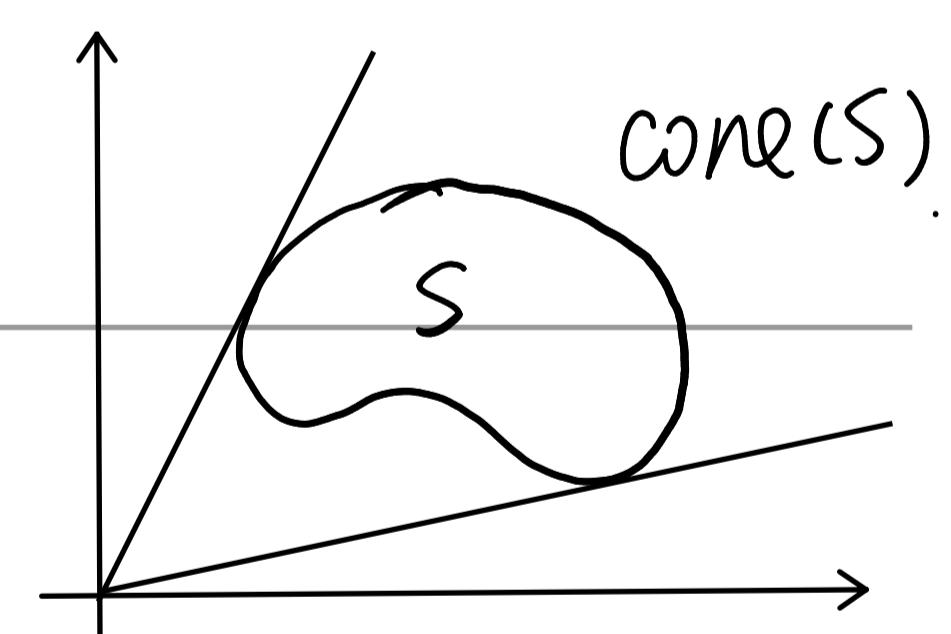


Application: Farkas' lemma. (LP duality).

Let $A \in \mathbb{R}^{n \times m}$, $b \in \mathbb{R}^n$. Then exactly one of the following is ϕ :

$$\textcircled{1} \quad \{x \in \mathbb{R}^m : Ax = b, x \geq \vec{0}\} \quad \textcircled{2} \quad \{y \in \mathbb{R}^n : A^T y \leq 0, b^T y > 0\}$$

Proof. Recall the conic combination $\theta_1 x_1 + \dots + \theta_k x_k$.



$$\text{cone}(S) \triangleq \left\{ \sum \theta_i x_i : \forall \theta_i \geq 0, x_i \in S \right\}$$

Let $A = (a_1, a_2, \dots, a_n)$. $\textcircled{1} = \phi \Rightarrow b \notin \text{cone}(A)$. closed, convex

By separating hyperplane (strict). $\exists y, z$. $u^T y + z < 0$ & $u^T y + z > 0$

$$\forall a_i, \forall \lambda_i \geq 0. \quad \lambda_i a_i^T y + z < 0 \Rightarrow a_i^T y + \frac{z}{\lambda_i} < 0 \Rightarrow a_i^T y \leq 0$$

$0 \in \text{cone}(A) \Rightarrow z < 0 \Rightarrow b^T y > 0$. so y is a desired one. \square