## Lecture 15. Newton's Method with Equality Constraints

### 15.1 Lagrangian function

Consider an optimization problem

$$
\begin{array}{rl}
\min _{\boldsymbol{x} \in \mathbb{R}^{n}} & f(\boldsymbol{x}) \\
\text { subject to } & g_{i}(\boldsymbol{x})=0, \quad \forall i=1,2, \ldots, m .
\end{array}
$$

We can reformulate the Lagrange condition in a compact form. We define the Lagrangian function as

$$
\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})=f(\boldsymbol{x})+\boldsymbol{\lambda}^{\top} \boldsymbol{g}(\boldsymbol{x})=f(\boldsymbol{x})+\lambda_{1} g_{1}(\boldsymbol{x})+\cdots+\lambda_{m} g_{m}(\boldsymbol{x}) .
$$

The Lagrange condition shows that if $\boldsymbol{x}^{*}$ is regular and local minimum, then there exists $\boldsymbol{\lambda}^{*}$ such that

$$
\nabla f\left(\boldsymbol{x}^{*}\right)+\left(\boldsymbol{\lambda}^{*}\right)^{\top} \boldsymbol{g}^{\prime}\left(\boldsymbol{x}^{*}\right)=\nabla f\left(\boldsymbol{x}^{*}\right)+\lambda_{1}^{*} \nabla g_{1}\left(\boldsymbol{x}^{*}\right)+\cdots+\lambda_{m}^{*} \nabla g_{m}\left(\boldsymbol{x}^{*}\right)=\mathbf{0} .
$$

Given the definition of Lagrangian, we can rewrite it as

$$
\nabla_{\boldsymbol{x}} \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0} .
$$

Note that

$$
\nabla_{\lambda} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})=\left(\frac{\partial \mathcal{L}}{\partial \lambda_{1}}, \ldots, \frac{\partial \mathcal{L}}{\partial \lambda_{m}}\right)^{\top}=\left(g_{1}(\boldsymbol{x}), \ldots, g_{m}(\boldsymbol{x})\right)^{\top}=\mathbf{0}
$$

if $\boldsymbol{x}$ is feasible. So we can simplify the Lagrange condition as $\nabla \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$.
Conversely, for convex optimization problems, the Lagrange condition (combining with the feasibility) is sufficient for optimal solutions, which is exactly $\nabla_{\boldsymbol{x}} \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$ and $\nabla_{\boldsymbol{\lambda}} \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$. Thus, we have the following sufficient and necessary condition for convex optimizations.

For any point $\boldsymbol{x}^{*} \in \mathbb{R}^{n}$, it is optimal for the convex optimization problem if and only if there exists $\boldsymbol{\lambda}^{*} \in \mathbb{R}^{m}$ such that

$$
\nabla \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}
$$

Note that $\mathcal{L}$ is not convex in general (see e.g., $f(x)=x^{2}, g(x)=x-5$, and thus $\mathcal{L}(x, \lambda)=x^{2}+\lambda(x-5)$ ), although $\mathcal{L}$ is a convex function of $\boldsymbol{x}$ for any fixed $\boldsymbol{\lambda}$. So $\nabla \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$ does not imply $\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)$ is a minimum point of $\mathcal{L}$. In fact, it is a saddle point in a sense. For feasible $\boldsymbol{x}$, i.e., $g_{1}(\boldsymbol{x})=\cdots=g_{m}(\boldsymbol{x})=0$, it is easy to see that $f(\boldsymbol{x})=\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})$. For infeasible $\boldsymbol{x}$, there exists $g_{i}(\boldsymbol{x}) \neq 0$, so we have $\max _{\boldsymbol{\lambda}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}) \rightarrow \infty$. Therefore, we conclude that

$$
\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\min _{\boldsymbol{x}} \max _{\boldsymbol{\lambda}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})
$$

Moreover, for convex optimization problems, the other direction also holds, namely, we have the following proposition.

## Theorem

For the Lagrangian $\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})$ of a convex optimization problem, if $\nabla \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$, then

$$
f\left(\boldsymbol{x}^{*}\right)=\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\max _{\boldsymbol{\lambda}} \min _{\boldsymbol{x}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})
$$

## Proof

For the first equality, note that if $\nabla \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}, g_{1}\left(\boldsymbol{x}^{*}\right)=\cdots=g_{m}\left(\boldsymbol{x}^{*}\right)=0$.
Thus $f\left(\boldsymbol{x}^{*}\right)=\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)$.
Now we define $\tilde{\mathcal{L}}(\boldsymbol{\lambda})=\min _{\boldsymbol{x}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})$. Our goal is to show that
$\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\max _{\boldsymbol{\lambda}} \tilde{\mathcal{L}}(\boldsymbol{\lambda})$.
We first show that there exists $\tilde{\boldsymbol{\lambda}}$ such that $\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\tilde{\mathcal{L}}(\tilde{\boldsymbol{\lambda}})$. In particular, we show that $\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\tilde{\mathcal{L}}\left(\boldsymbol{\lambda}^{*}\right)$. Note that $\tilde{\mathcal{L}}\left(\boldsymbol{\lambda}^{*}\right)=\min _{\boldsymbol{x}} \mathcal{L}\left(\boldsymbol{x}, \boldsymbol{\lambda}^{*}\right)$. Since $\boldsymbol{\lambda}^{*}$ is fixed, $\mathcal{L}\left(\boldsymbol{x}, \boldsymbol{\lambda}^{*}\right)$ is a convex function of $\boldsymbol{x}$. By the first order condition for convexity, $\mathcal{L}\left(\boldsymbol{x}, \boldsymbol{\lambda}^{*}\right)$ achieves the minimum value if and only if $\nabla_{\boldsymbol{x}} \mathcal{L}\left(\boldsymbol{x}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$ . Since $\nabla_{\boldsymbol{x}^{*}} \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)=\mathbf{0}$, we have $\tilde{\mathcal{L}}\left(\boldsymbol{\lambda}^{*}\right)=\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)$.
Next, we show that for all $\boldsymbol{\lambda} \in \mathbb{R}^{m}, \tilde{\mathcal{L}}(\boldsymbol{\lambda}) \leq \mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)$. Denote by $\Omega$ the feasible set $\left\{\boldsymbol{x} \mid g_{1}(\boldsymbol{x})=g_{2}(\boldsymbol{x})=\cdots=g_{m}(\boldsymbol{x})=0\right\}$. Then for all $\boldsymbol{x} \in \Omega$ and all $\boldsymbol{\lambda}$,

$$
\begin{aligned}
& f(\boldsymbol{x})=\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}) . \text { So for all } \boldsymbol{\lambda} \in \mathbb{R}^{m}, \\
& \quad \tilde{\mathcal{L}}(\boldsymbol{\lambda})=\min _{\boldsymbol{x}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}) \leq \min _{\boldsymbol{x} \in \Omega} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})=\min _{\boldsymbol{x} \in \Omega} f(\boldsymbol{x})=f\left(\boldsymbol{x}^{*}\right)=\mathcal{L}\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right) .
\end{aligned}
$$

Later, we will revisit this theorem in the context of general convex optimization with inequality constraints.

### 15.2 Karush-Kuhn-Tucker System

We now consider how to solve convex optimization problems with equality constraints. Recall the problems with no constraints, where we approximate the objective function by a quadratic function and minimize the quadratic function. We would like to apply the same idea. So the first step is to minimize quadratic functions with equality constraints.

Given $\boldsymbol{Q} \in \mathbb{R}^{n \times n} \succeq \mathbf{0}, \boldsymbol{A} \in \mathbb{R}^{m \times n}, \boldsymbol{w} \in \mathbb{R}^{n}$ and $\boldsymbol{b} \in \mathbb{R}^{m}$, consider the following quadratic problem with equality constrains

$$
\begin{aligned}
\min _{\boldsymbol{x} \in \mathbf{R}^{n}} & \frac{1}{2} \boldsymbol{x}^{\top} \boldsymbol{Q} \boldsymbol{x}+\boldsymbol{w}^{\top} \boldsymbol{x} \\
\text { subject to } & \boldsymbol{A} \boldsymbol{x}=\boldsymbol{b}
\end{aligned}
$$

We can solve this problem by the Lagrange multiplier method. The Lagrangian function is

$$
\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})=\frac{1}{2} \boldsymbol{x}^{\top} \boldsymbol{Q} \boldsymbol{x}+\boldsymbol{w}^{\top} \boldsymbol{x}+\boldsymbol{\lambda}^{\top}(\boldsymbol{A} \boldsymbol{x}-\boldsymbol{b})
$$

and the Lagrange condition is

$$
\left\{\begin{array}{l}
\nabla_{\boldsymbol{x}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})=\boldsymbol{Q} \boldsymbol{x}+\boldsymbol{w}+\boldsymbol{A}^{\top} \boldsymbol{\lambda}=\mathbf{0} \\
\nabla_{\boldsymbol{\lambda}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda})=\boldsymbol{A} \boldsymbol{x}-\boldsymbol{b}=\mathbf{0}
\end{array}\right.
$$

We know that $\boldsymbol{x}^{*}$ is a global minimum if and only $\left(\boldsymbol{x}^{*}, \boldsymbol{\lambda}^{*}\right)$ is a solution to the above system of equations for some $\boldsymbol{\lambda}^{*}$. The above system of equations can be expressed in the following matrix form:

$$
\underbrace{\left(\begin{array}{cc}
\boldsymbol{Q} & \boldsymbol{A}^{\top} \\
\boldsymbol{A} & \mathbf{0}
\end{array}\right)\binom{\boldsymbol{x}}{\boldsymbol{\lambda}}=\binom{-\boldsymbol{w}}{\boldsymbol{b}}}_{\text {KKT system }}
$$

which is called the KKT system and the coefficient matrix $\left(\begin{array}{cc}\boldsymbol{Q} & \boldsymbol{A}^{\top} \\ \boldsymbol{A} & \mathbf{0}\end{array}\right)$ is called the KKT matrix.

## Block Gaussian elimination

If we further assume $\boldsymbol{Q}$ is invertible (nonsingular), we can solve the KKT system by Gaussian elimination. Because $\boldsymbol{Q} \succeq \mathbf{0}$, it is equivalent to $\boldsymbol{Q} \succ \mathbf{0}$.

Left multiplying $\boldsymbol{A} \boldsymbol{Q}^{-1}$ to the first row and subtracting the second row, we obtain that

$$
\left(\begin{array}{ll}
0 & \boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}
\end{array}\right)\binom{\boldsymbol{x}}{\boldsymbol{\lambda}}=-\boldsymbol{b}-\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{w}
$$

So

$$
\boldsymbol{\lambda}=-\left(\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}\right)^{-1}\left(\boldsymbol{b}+\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{w}\right)
$$

Plugging it into the first row of the original KKT system

$$
\boldsymbol{Q} \boldsymbol{x}+\boldsymbol{A}^{\top} \boldsymbol{\lambda}=-\boldsymbol{w},
$$

we have

$$
\boldsymbol{x}=-\boldsymbol{Q}^{-1} \boldsymbol{w}+\boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}\left(\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}\right)^{-1}\left(\boldsymbol{b}+\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{w}\right)
$$

## Remark

Both $\boldsymbol{Q}^{-1}$ and $\left(\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}\right)^{-1}$ must exist in the above calculation. If we assume $\boldsymbol{Q} \succ 0$ and $\operatorname{rank}(\boldsymbol{A})=m$ ( $\boldsymbol{A}$ has full row rank), then both of them exist. First, $\boldsymbol{Q} \succ \mathbf{0}$ implies that $\boldsymbol{Q}$ is invertible. Next, if $\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}$ is not invertible, there exists $\boldsymbol{v} \in \mathbb{R}^{m} \neq \mathbf{0}$ such that $\left(\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}\right) \boldsymbol{v}=\mathbf{0}$. If we left multiple $\boldsymbol{v}^{\top}$ to both sides, we can get

$$
\boldsymbol{v}^{\top}\left(\boldsymbol{A} \boldsymbol{Q}^{-1} \boldsymbol{A}^{\top}\right) \boldsymbol{v}=\left(\boldsymbol{A}^{\top} \boldsymbol{v}\right)^{\top} \boldsymbol{Q}^{-1}\left(\boldsymbol{A}^{\top} \boldsymbol{v}\right)=0 .
$$

Since $\operatorname{rank}(\boldsymbol{A})=m$, its rows are linear independent. So $\boldsymbol{A}^{\top} \boldsymbol{v} \neq \mathbf{0}$, and $\left(\boldsymbol{A}^{\top} \boldsymbol{v}\right)^{\top} \boldsymbol{Q}^{-1}\left(\boldsymbol{A}^{\top} \boldsymbol{v}\right)=0$ contradicts to the fact that $\boldsymbol{Q}^{-1} \succ \mathbf{0}$.
The hypothesis of $\operatorname{rank}(\boldsymbol{A})=m$ is reasonable. Otherwise either the problem is infeasible (e.g., constraints are $x_{1}+x_{2}=2$ and $2 x_{1}+2 x_{2}=3$ ), or there are redundant constraints (e.g., constraints are $x_{1}+x_{2}=2$ and $2 x_{1}+2 x_{2}=4$ ).

But the assumption of $\boldsymbol{Q} \succ 0$ is strong. Actually, we only need the KKT matrix to be invertible.

## Nonsingularity of KKT matrices

Let $f(\boldsymbol{x})=\frac{1}{2} \boldsymbol{x}^{\top} \boldsymbol{Q} \boldsymbol{x}+\boldsymbol{w}^{\boldsymbol{\top}} \boldsymbol{x}$ where $\boldsymbol{Q} \succeq \mathbf{0}$. If there is no constraints, when does $f$ attains it minimum value? The gradient is $\boldsymbol{Q} \boldsymbol{x}+\boldsymbol{w}$. So there exists an optimal solution if and only if the equation $\boldsymbol{Q} \boldsymbol{x}+\boldsymbol{w}=\mathbf{0}$ has solutions, namely, $\boldsymbol{w}$ is in the row space of $\boldsymbol{Q}$. Recall the lemma $\operatorname{im}\left(\boldsymbol{A}^{\top}\right)=\operatorname{ker}(\boldsymbol{A})^{\perp}$. It is further equivalent to $\boldsymbol{w} \perp \operatorname{ker}(\boldsymbol{Q})$, i.e.,

$$
\forall \boldsymbol{v} \in \mathbb{R}^{n}, \quad \boldsymbol{Q} \boldsymbol{v}=\mathbf{0} \Longrightarrow\langle\boldsymbol{w}, \boldsymbol{v}\rangle=0
$$

There are two cases:

1. if $\boldsymbol{Q}$ is invertible / nonsingular, there exists a unique solution $\boldsymbol{x}$;
2. if $\boldsymbol{Q}$ is not invertible, there are infinite many solutions.

Now let $\boldsymbol{K}$ be the KKT matrix. Clearly the KKT system is solvable if and only if $\binom{-\boldsymbol{w}}{\boldsymbol{b}} \perp \operatorname{ker}(\boldsymbol{K})$.

We first prove the following useful proposition.

## Lemma

Suppose $Q \succeq 0$, then for any vector $x, Q x=\mathbf{0}$ if and only if $x^{\top} Q x=0$.

## Proof

- Necessity: Trivial.
- Sufficiency: Since $Q \succeq 0$, we can do eigendecomposition to $Q$.

$$
Q=U \Lambda U^{\top}=\left[\xi_{1} u_{1}, \ldots, \xi_{n} u_{n}\right] \cdot\left(\begin{array}{c}
u_{1}^{\top} \\
u_{2}^{\top} \\
\ldots \\
u_{n}^{\top}
\end{array}\right)=\sum_{i=1}^{n} \xi_{i} u_{i} u_{i}^{\top},
$$

where $\xi_{i} \geq 0$ is the $i$-th eigenvalue and $u_{i}$ is the $i$-th eigenvector of $Q$.

Then we have

$$
x^{\top} Q x=\sum_{i=1}^{n} \xi_{i}\left(u_{i}^{\top} x\right)^{2}=0,
$$

which means either $\xi_{i}=0$ or $u_{i}^{\top} x=0$. Hence,

$$
Q x=\sum_{i=1}^{n} u_{i}\left(\xi_{i}\left(u_{i}^{\top} x\right)\right)=\mathbf{0} .
$$

Now we can show that $\operatorname{ker}(\boldsymbol{K})=\left\{(\boldsymbol{v}, \mathbf{0})^{\top} \mid \boldsymbol{v} \in \operatorname{ker}(\boldsymbol{Q}) \cap \operatorname{ker}(\boldsymbol{A})\right\}$. Note that if $\boldsymbol{v} \in \operatorname{ker}(\boldsymbol{Q}) \cap \operatorname{ker}(\boldsymbol{A})$, then

$$
\left(\begin{array}{cc}
\boldsymbol{Q} & \boldsymbol{A}^{\top} \\
\boldsymbol{A} & \mathbf{0}
\end{array}\right)\binom{\boldsymbol{v}}{\mathbf{0}}=\binom{-\boldsymbol{Q} \boldsymbol{v}}{\boldsymbol{A} \boldsymbol{v}}=\mathbf{0} .
$$

Conversely, we claim that if

$$
\left(\begin{array}{cc}
Q & A^{\boldsymbol{\top}} \\
A & \mathbf{0}
\end{array}\right)\binom{u}{v}=\mathbf{0},
$$

then $u \in \operatorname{ker}(Q) \cap \operatorname{ker}(A)$ and $v=\mathbf{0}$. This linear equation system is equivalent to $Q u+A^{\top} v=\mathbf{0}$ and $A u=\mathbf{0}$. Note that by left multiplying $u^{\top}$, we obtain that

$$
u^{\top} Q u=u^{\top}\left(-A^{\top} v\right)=-(A u)^{\top} v=0
$$

So $Q u=\mathbf{0}$ by the above lemma. Moreover, $A^{\top} v=-Q u=\mathbf{0}$, contradicts $\operatorname{rank}(A)=m$ if $v \neq \mathbf{0}$.

Then we focus on the case where $\boldsymbol{K}$ is invertible / nonsingular.

## Theorem

KKT matrix is invertible (nonsingular) is equivalent to each one of the followings:

1. $\operatorname{ker}(A) \cap \operatorname{ker}(Q)=\{\mathbf{0}\}$, or
2. if $A x=0$ and $x \neq 0$, then $x^{\top} Q x>0$, or
3. $\forall F \in \mathbb{R}^{n \times(n-m)}$, if $\operatorname{im}(F) \triangleq\left\{F v: v \in \mathbb{R}^{n-m}\right\}=\operatorname{ker}(A)$, then $F^{\top} Q F \succ 0$.

- "invertible $\Longrightarrow 1$ ". If there exists $x \neq \mathbf{0}$ such that $x \in \operatorname{ker}(Q) \cap \operatorname{ker}(A)$, we have

$$
\left(\begin{array}{cc}
Q & A^{\top} \\
A & \mathbf{0}
\end{array}\right) \cdot\binom{x}{\mathbf{0}}=\mathbf{0},
$$

which means the KKT matrix is not invertible.

- " $1 \Longrightarrow 2$ ". If $A x=0$ and $x \neq 0$, then $x \in \operatorname{ker}(A) \backslash\{\mathbf{0}\}$. So $Q x \neq 0$. Thus $x^{\top} Q x \neq 0$ by the above lemma.
- " $2 \Longrightarrow$ invertible". By the above lemma, item 2 is equivalent to if $A x=0$ and $x \neq 0$ then $Q x \neq \mathbf{0}$. So $\operatorname{ker}(A) \cap \operatorname{ker}(Q)=\{\mathbf{0}\}$. Thus $\operatorname{ker}(\boldsymbol{K})=\{\mathbf{0}\}$, which implies that $\boldsymbol{K}$ is invertible.
- " $2 \Longrightarrow 3$ ". If $\operatorname{im}(F)=\operatorname{ker}(A)$, for all $x \neq \mathbf{0}, A F x=0$. Note that $\operatorname{dimim}(F)=\operatorname{dim} \operatorname{rank}(A)=n-m$. So $\operatorname{rank}(F)=n-m$. Thus $F x \neq \mathbf{0}$ by linear independence. If $F x \neq \mathbf{0}$ then $(F x)^{\top} Q(F x)=x^{\top}\left(F^{\top} Q F\right) x>0$. So $F^{\top} Q F \succ 0$.
- " $3 \Longrightarrow 2$ ". If $\operatorname{im}(F)=\operatorname{ker}(A)$, for all $x \neq \mathbf{0}$ that $A x=\mathbf{0}$, there exists $y \neq \mathbf{0}$ such that $x=F y$. Since $F^{\top} Q F \succ 0, x^{\top} Q x=y^{\top} F^{\top} Q F y>0$.


### 15.3 Newton's method

Now we consider how to solve general convex optimization with equality constraints.

Recall Newton's method. Given $x_{k}$, we do Taylor Expansion of $f(x)$ at $x_{k}$ :

$$
f(x) \approx \tilde{f}(x) \triangleq f\left(x_{k}\right)+\nabla f\left(x_{k}\right)^{\top}\left(x-x_{k}\right)+\frac{1}{2}\left(x-x_{k}\right)^{\top} \nabla^{2} f\left(x_{k}\right)\left(x-x_{k}\right)
$$

- If there is no constraints, we use $x_{k+1}=\arg \min \tilde{f}(x)$ to approximate the minimum point of $f$.
- If there are constraints $A x-b=0$, we may use $x_{k+1}=\underset{A x-b=0}{\arg \min } \tilde{f}(x)$ to approximate the minimum point of $f$ under constraints $A x-b=\mathbf{0}$.

Let $d=x-x_{k}$. We have

$$
\tilde{f}(x)=f\left(x_{k}\right)+\nabla f\left(x_{k}\right)^{\top} d+\frac{1}{2} d^{\top} \nabla^{2} f\left(x_{k}\right) d
$$

and the constraints become $A d=\mathbf{0}$ since

$$
0=A\left(x_{k}+d\right)-b=A x_{k}+A d-b .
$$

If $x_{k}$ is a feasible solution, we have $A x_{k}-b=\mathbf{0}$, hence $A d=\mathbf{0}$.
Now for $d$, the problem becomes

$$
\underset{A d=0}{\arg \min } \tilde{f}(x)=f\left(x_{k}\right)+\nabla f\left(x_{k}\right)^{\top} d+\frac{1}{2} d^{\top} \nabla^{2} f\left(x_{k}\right) d
$$

and finally

$$
x_{k+1}=\underset{A x-b=\mathbf{0}}{\arg \min } \tilde{f}(x)=x_{k}+\underset{A d=\mathbf{0}}{\arg \min } f\left(x_{k}\right)+\nabla f\left(x_{k}\right)^{\top} d+\frac{1}{2} d^{\top} \nabla^{2} f\left(x_{k}\right) d .
$$

The KKT system for optimizing $d$ is

$$
\left(\begin{array}{cc}
\nabla^{2} f\left(x_{k}\right) & A^{\top} \\
A & \mathbf{0}
\end{array}\right)\binom{d}{\lambda}=\binom{-\nabla f\left(x_{k}\right)}{0}
$$

```
ALGORITHM 1: Formal Definition of Newton's Method
Let \(x_{0}\) be an arbitrary feasible solution;
while \(d^{T} \nabla^{2} f\left(x_{k}\right) d \leq \delta\) do
    Compute \(d\) by solving KKT system;
    \(x_{k+1} \leftarrow x_{k}+d\);
return \(\left\{x_{k}\right\}\)
```

Note that we use $d^{\top} \nabla^{2} f\left(x_{k}\right) d \leq \delta$ as stopping criteria instead of $\left\|\nabla f\left(x_{k}\right)\right\|<\delta$. This is because, when the algorithm should terminate, its gradient is not zero (it needs to follow the Lagrange condition).

Next, let us see some examples of using Newton's method.

## Example 1

$$
\begin{array}{cl}
\min & f\left(x_{1}, x_{2}\right)=x_{1}^{2}+x_{2}^{2} \\
\text { s.t. } & x_{1}+x_{2}=1
\end{array}
$$

Start from $\binom{1}{0}$. The KKT system at this point is

$$
\left(\begin{array}{lll}
2 & 0 & 1 \\
0 & 2 & 1 \\
1 & 1 & 0
\end{array}\right)\left(\begin{array}{c}
d_{1} \\
d_{2} \\
\lambda
\end{array}\right)=\left(\begin{array}{c}
-2 \\
0 \\
0
\end{array}\right)
$$

Its solution is $\left(d_{1}, d_{2}, \lambda\right)=\left(-\frac{1}{2}, \frac{1}{2},-1\right)$. The next point $x_{1}$ is exactly the optimal solution $\left(\frac{1}{2}, \frac{1}{2}\right)$.

## Example 2

$$
\begin{array}{cl}
\min & f\left(x_{1}, x_{2}\right)=x_{1}^{2} \\
\text { s.t. } & x_{1}+2 x_{2}=b
\end{array}
$$

Start from $\binom{b}{0}$. The KKT system at this point is

$$
\left(\begin{array}{lll}
2 & 0 & 1 \\
0 & 0 & 2 \\
1 & 2 & 0
\end{array}\right)\left(\begin{array}{c}
d_{1} \\
d_{2} \\
\lambda
\end{array}\right)=\left(\begin{array}{c}
-2 b \\
0 \\
0
\end{array}\right) .
$$

Its solution is $\left(d_{1}, d_{2}, \lambda\right)=\left(-b, \frac{b}{2}, 0\right)$. The next point $x_{1}$ is also exactly the optimal solution ( $0, \frac{b}{2}$ ).
In this example, although $\nabla^{2} f$ is not invertible, the KKT matrix is still invertible.
We may check by the criterion introduced above, since
$\operatorname{ker}\left(\nabla^{2} f\right)=\left\{s \cdot\binom{0}{1}: s \in \mathbb{R}\right\}$ and $\operatorname{ker}(A)=\left\{t \cdot\binom{-2}{1}: t \in \mathbb{R}\right\}$, which satisfy the previous condition of $\operatorname{ker}(A) \cap \operatorname{ker}\left(\nabla^{2} f\right)=\{0\}$.

## Example 3

$$
\begin{array}{cl}
\min & f\left(x_{1}, x_{2}\right)=e^{x_{1}^{2}+x_{2}^{2}} \\
\text { s.t. } & x_{1}+x_{2}=1
\end{array}
$$

Start from $\binom{1}{0}$. The KKT system at this point is

$$
\left(\begin{array}{ccc}
6 e & 0 & 1 \\
0 & 2 e & 1 \\
1 & 1 & 0
\end{array}\right)\left(\begin{array}{c}
d_{1} \\
d_{2} \\
\lambda
\end{array}\right)=\left(\begin{array}{c}
-2 e \\
0 \\
0
\end{array}\right) .
$$

Its solution is $\left(d_{1}, d_{2}, \lambda\right)=\left(-\frac{1}{4}, \frac{1}{4},-\frac{e}{2}\right)$. The next point $x_{1}$ is $\left(\frac{3}{4}, \frac{1}{4}\right)$. It is not the optimal solution, but it is also in the right direction.

### 15.4 Correctness and convergence

We will demonstrate the correctness of Newton's method via the following two propositions:

1. Each time we choose a descending direction: $\nabla f\left(x_{k}\right)^{\top} d \leq 0$.
2. The stopping criteria is correct: if $d^{\top} \nabla^{2} f\left(x_{k}\right) d=0$ then $x_{k}$ is optimal.

## Proposition 1

$\nabla f\left(x_{k}\right)^{\top} d \leq 0$.

## Proof

The KKT system can be unfolded as follows:

$$
\left\{\begin{array}{l}
\nabla^{2} f\left(x_{k}\right) d+A^{\top} \lambda=-\nabla f\left(x_{k}\right) \\
A d=0
\end{array}\right.
$$

Use $d^{\top}$ to left multiple the first line, we can get

$$
d^{\top} \nabla^{2} f\left(x_{k}\right) d+\underbrace{d^{\top} A^{\top} \lambda}_{=0}=-d^{\top} \nabla f\left(x_{k}\right),
$$

which yields

$$
d^{\top} \nabla f\left(x_{k}\right)=-d^{\top} \nabla^{2} f\left(x_{k}\right) d \leq 0 .
$$

## Proposition 2

If $d^{\top} \nabla^{2} f\left(x_{k}\right) d=0$ then $x_{k}$ is optimal.

## Proof

First, $d^{\top} \nabla^{2} f\left(x_{k}\right) d=0 \Leftrightarrow \nabla^{2} f\left(x_{k}\right) d=\mathbf{0}$ (due to the lemma proved in Section 15.2). Then by the first equality in KKT system,

$$
\underbrace{\nabla^{2} f\left(x_{k}\right) d}_{=0}+A^{\top} \lambda=-\nabla f\left(x_{k}\right) \Longrightarrow \nabla f\left(x_{k}\right)+A^{\top} \lambda=0 \Longrightarrow x_{k} \text { is optimal. }
$$

Now we analyzed the convergence of Newton's method with equality constraints. In fact, we can convert equality constrained problems to problems without
constraint. For example, the following problem

$$
\begin{array}{cl}
\min & f\left(x_{1}, x_{2}\right) \\
\text { s.t. } & x_{1}+x_{2}=1
\end{array}
$$

is equivalent to $\min f\left(x_{1}, 1-x_{1}\right)$.
In general, for an equality-constrainted problem, assume its feasible set is $\Omega=\{x \mid A x=b\}$. It can be rewritten as $\Omega=\tilde{x}+\left\{F z \mid z \in \mathbb{R}^{n-m}\right\}$ for some $\tilde{x} \in \Omega$ and $F \in \mathbb{R}^{n \times(n-m)}$, since $\Omega$ is an affine set.
Then the original problem is equivalent to the following one:

$$
\min _{z \in \mathbb{R}^{n-m}} g(z) \triangleq f(\tilde{x}+F z)
$$

Applying Newton's method to $g(z)$, we will get

$$
\begin{aligned}
z_{k+1} & =z_{k}-\left(\nabla^{2} g\left(z_{k}\right)\right)^{-1} \nabla g\left(z_{k}\right) \\
& =z_{k}-\left(\nabla\left(F^{\top} \nabla f\left(\tilde{x}+F z_{k}\right)\right)\right)^{-1} \nabla g\left(z_{k}\right) \\
& =z_{k}-\left(F^{\top} \nabla^{2} f\left(\tilde{x}+F z_{k}\right) F\right)^{-1} \nabla\left(F^{\top} \nabla f\left(\tilde{x}+F z_{k}\right)\right)
\end{aligned}
$$

Assuming $x_{0}=\tilde{x}+F z_{0}$, we have the following proposition, which shows that the Newton's method is affinely invariant.

## Lemma

For all $k \geq 1, x_{k}=\tilde{x}+F z_{k}$.

## Proof

By induction, assume $x_{k}=\tilde{x}+F z_{k}$. Let $d_{x_{k}}=x_{k+1}-x_{k}$ and $d_{z_{k}}=z_{k+1}-z_{k}$. For all $v \in \mathbb{R}^{n}$, if $A v=\mathbf{0}$ then $\tilde{x}+v \in \Omega$, which implies $v \in \operatorname{im}(F)$, and vice versa. So $\operatorname{im}(F)$ is exactly $\operatorname{ker}(A)$.
Note that $A d_{x_{k}}=0$ by Newton's method. Thus there exists $u \in \mathbb{R}^{n-m}$ such that $d_{x_{k}}=F u$.
Moreover, by the first equalty of ККТ system, we know that

$$
\nabla^{2} f\left(x_{k}\right) d_{x_{k}}+A^{\top} \lambda=-\nabla f\left(x_{k}\right) \Rightarrow \nabla^{2} f\left(x_{k}\right) F u+A^{\top} \lambda=-\nabla f\left(x_{k}\right)
$$

Left multiply $F^{\top}$ on both sides,

$$
F^{\top} \nabla^{2} f\left(x_{k}\right) F u+F^{\top} A^{\top} \lambda=-F^{\top} \nabla f\left(x_{k}\right)
$$

Since $\operatorname{im}(F)$ is $\operatorname{ker}(A)$, we have $A F z=0$ for all $z$, which implies that $A F=\mathbf{0}$.

Then, we have

$$
F^{\top} \nabla^{2} f\left(x_{k}\right) F u=-F^{\top} \nabla f\left(x_{k}\right),
$$

which implies that

$$
u=-\left(F^{\top} \nabla^{2} f\left(x_{k}\right) F\right)^{-1}\left(F^{\top} \nabla f\left(x_{k}\right)\right)=-\left(\nabla^{2} g\left(z_{k}\right)\right)^{-1} \nabla g\left(z_{k}\right)=d_{z_{k}} .
$$

Hence, $d_{x_{k}}=F u=F d_{z_{k}}$ and

$$
x_{k+1}=x_{k}+d_{x_{k}}=x_{k}+F d_{z_{k}}=\tilde{x}+F z_{k+1} .
$$

Therefore, the convergence of $\left\{z_{k}\right\}$ can lead to the convergence of $\left\{x_{k}\right\}$.

